# A Multi-tool for your Quantum Algorithmic Toolbox



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Based on work with

- Chloe Ye, Da-Yeon Koh, Noel Anderson, Jay-U Chung arXiv:2012.01276
- Teal Witter arXiv:2010.02324 (WADS 2021)
- Kai DeLorenzo, Teal Witter, arXiv:1904.05995 (TQC 2019)
- Michael Jarret, Stacey Jeffery, Alvaro Piedrafita, arXiv:1804.10591 (ESA 2018)
- Stacey Jeffery: arXiv: 1704.00765 (Quantum vol 1 p 26)
- Bohua Zhan, Avinatan Hassidim, arXiv:1101.0796 (ITCS 2012)

Quantum computers use quantum particles – atoms, photons, phonons, electrons, etc. – to compute.

Famously fast for:

- factoring
- simulating quantum chemistry

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• Specialized Problems

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- Specialized Problems
- Sophisticated design/analysis

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- Specialized
   Problems
- Sophisticated design/analysis
- Everyday Problems?
- Accessible design/analysis?

#### Year 2027



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|                |    |   |   |   |         |  |  |

#### Year 2027



#### # else:

return -1.0\*(abs(test1)+abs(test2)) #

#### def sim\_mat\_create(vec):

return np.array([[1.0+vec[0],vec[1],0],[vec[2],1.0+vec[3],0 #return np.array([[1.0+vec[0],vec[1],0],[0.0,1.0,0],[0.0,0.0]

#### def sim\_mat\_create\_full(vec):

- return np.array([[1.0,0.0,0.0,0.0],[0.0,1.0+vec[0],vec[
- def sim\_transform(sim\_mat,target\_map): sim\_mat\_inv=inv(sim\_mat) return dot(dot(sim\_mat,target\_map),sim\_mat\_inv)
- def test\_transform(vec,target\_map): return SimpleCPTest(sim\_transform(sim\_mat\_create(vec),targe

#### def maxfunc(vec,matx,matz):

sim\_mat=np.array([[1.0+vec[0],vec[1],0],[vec[2],1.0+vec[3], #sim\_mat=np.array([[1.0+vec[0],vec[1],0],[0.0,1.0,0],[0.0,0 sim\_mat\_inv=inv(sim\_mat) m1=dot(dot(sim\_mat,matx),sim\_mat\_inv) m2=dot(dot(sim\_mat,matz),sim\_mat\_inv) value = Two\_CP\_matrices(m1,m2) return -1\*value

#### #li=np.array([.0,.0001,.0,.0])

li=np.array([0,0,0,0]) #li=np.array([0,0.0]) res = minimize(maxfunc, li, (gxsub, gzsub), method='nelder-mead output=res.x print(test\_transform(output,gzsub)) print(test\_transform(output,gxsub)) print(output)

Optimization terminated successfully. Current function value: -0.000123 Iterations: 233 Function evaluations: 415 0 000123103176507



#### Year 2027



!lse: return -1.0\*(abs(test1)+abs(test2))

im\_mat\_create(vec):
sturn np.array([[1.0+vec[0],vec[1],0],[vec[2],1.0+vec[3],0])
teturn np.array([[1.0+vec[0],vec[1],0],[0.0,1.0,0],[0.0,0.4])

im\_mat\_create\_full(vec):
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im\_transform(sim\_mat,target\_map): im\_mat\_inv=inv(sim\_mat) sturn dot(dot(sim\_mat,target\_map),sim\_mat\_inv)

st\_transform(vec,target\_map):
turn SimpleCPTest(sim\_transform(sim\_mat\_create(vec),targe)

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#### im\_mat\_create\_full(vec): return np.array([[1.0,0.0,0.0,0.0],[0.0,1.0+vec[0],vec[

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#### Row With all 1's? 011 ... 00001 011 ... 10111 111 ... 11111 010 ... 01001 Path? Cycle? t

**Our Strategy** 

Reduce our problem to path detection



Reduce our problem to path detection

### Outline

- 1. Path Detection (set-up)
- 2. Reductions
- 3. Quantum Path Detection Algorithm (complexity)



Reduce our problem to path detection

Is there a path from s to t?



Bit String Input *x*:



"Skeleton graph" G

Is there a path from s to t?



Bit String Input *x*:



Is there a path from s to t?



Bit String Input *x*:



Is there a path from s to t?



Bit String Input *x*:



Catch:

• Bit string initially hidden







Query in superposition

$$\sum_{i} |i\rangle |0\rangle = O_{x} \qquad \Rightarrow \sum_{i} |i\rangle |x_{i}\rangle$$

(states not normalized)



Goal: Figure out if path while using the oracle as few times as possible (while revealing as few bits as possible) -> minimize **Query Complexity** 

For our algorithms, runtime scales like query complexity, times the time it takes to do one step of a quantum walk on the skeleton graph (for most graphs with structure: log time)

# Reduction to Path Detection







Reduce our problem to path detection

Problem: Is there a cycle?





Yes

No

Problem: Is there a cycle?



Unknown bit string  $x_1 x_2 x_3 x_4 x_5$ 

 $x_i = 1 \leftrightarrow \text{edge } i \text{ is present}$ 

Problem: Is there a cycle?

Subproblem: Is there a cycle through edge 1?



There is a cycle through Edge 1 iff

- Edge 1 is present 👡
- Path between the endpoints of Edge 1
   AND not using Edge 1

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Or

Problem: Is there a cycle?



There is a cycle if

- Cycle through edge 1
- Cycle through edge 2
- •
- Cycle through edge n

Problem: Is there a cycle?

2 5 3 4

There is a cycle if

- Cycle through edge 1
- Cycle through edge 2
- Cycle through edge n



Is there a path from s to t?



Bit String Input *x*:





G(x)

#### Formula Evaluation to Path Detection



[Nisan, Ta-Shma, '95]

#### Quantum Path Detection Reductions

- Read-once Boolean formulas [Jeffery, **K**, '17]
- Total connectivity [Jarret, Jeffery, K, Piedrafita, '18]
- Cycle detection [Delorenzo, K, Witter, '19]
- Even length cycle detection [Delorenzo, K, Witter, '19]
- Bipartiteness [Delorenzo, K, Witter, '19]
- Maximum matching (**K**, Witter, '21)
- 2-player game evaluation [Zhan, Hassidim, K, '12]
- Directed st-connectivity (Beigi, Taghavi `19)
- Directed smallest cycle (Beigi, Taghavi `19)
- Topological sort (Beigi, Taghavi `19)
- Connected components (Beigi, Taghavi `19)
- Strongly connected components (Beigi, Taghavi `19)
- k-cycle at vertex v (Beigi, Taghavi `19)
- st-connectivity (Reichardt, Belovs `12)
- Maximum bipartite matching (Lin, Lin `16; Beigi and Taghavi `19)



Reduce our problem to path detection



Reduce our problem to path detection

# Quantum Path Detection Algorithm Complexity

Space Complexity: O(log(# edges in skeleton graph))

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Space Complexity: O(log(# edges in skeleton graph))

Query Complexity: Depends on

- effective resistance
- effective capacitance

# Effective Resistance



# Effective Resistance

R(G(x)) is effective resistance of this circuit



# Effective Capacitance

G(x)



# Effective Capacitance



# Quantum Path Detection Algorithm Complexity

Quantum Query Complexity of Path Detection:

• If path in G(x):

$$\tilde{O}\left(\sqrt{R(G(x)) \times \max_{\substack{y:G(y)\\has no path}} C(G(y))}\right)$$

[Belovs, Reichardt '12; JJKP, `18; ACKKY '20]

# Quantum Path Detection Algorithm Complexity

Quantum Query Complexity of Path Detection:

• If path in G(x):



[Belovs, Reichardt '12; JJKP, `18; ACKKY '20]

Search: Is there a 1 in  $x = x_1 x_2 \dots x_n$ ?

Search: Is there a 1 in  $x = x_1 x_2 \dots x_n$ ?

Is there a 1 at bit 1? OR Is there a 1 at bit 2? OR Is there a 1 at bit 3? OR

...

Search: Is there a 1 in  $x = x_1 x_2 \dots x_n$ ?



No path:

• C(G(x)) = n



Path:

• R(G(x)) = 1/k (k is number of bits with value 1)



### Quantum Search Algorithm Complexity

Quantum Query Complexity of Path Detection:

• If path in G(x):



[Belovs, Reichardt '12; JJKP, `18; ACKKY '20]

### Quantum Search Algorithm Complexity

Quantum Query Complexity of Path Detection:

• If path in G(x):



Not new result (variant of Grover's Algorithm), but you designed it! (Within log factors of optimal)

# Performance

\*There are alternative optimal algorithms for some problems

- Read-once Boolean formulas (query optimal) [Jeffery, **K**, '17]
- Total connectivity (query optimal) [Jarret, Jeffery, **K**, Piedrafita, '18]
- Cycle detection (query optimal) [Delorenzo, K, Witter, '19]
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# Path Detection Algorithm:

$$U = R_1 R_2$$

Reflection that encodes quantum walk on the skeleton graph (Often  $O(\log n)$  time, no queries) Reflection that encodes where edges are in graph (Takes O(1) time and O(1) queries)

# Path Detection Algorithm:

Basic Algorithm (no speed up for easy instances):

Do eigenvalue estimation of U for some initial state with high enough precision

- If initial state has eigenvalue 1 -> decide path
- Otherwise -> decide no path

[Belovs, Reichardt '12]



Reduce our problem to path detection

# Future/Current Work

• Recent work (with Stacey Jeffery and Alvaro Piedrafita): algorithm to sample an edge in the flow

# Effective Resistance

#### R(G(x)) is effective resistance of this circuit



# Future/Current Work

- Recent work (with Stacey Jeffery and Alvaro Piedrafita): algorithm to sample an edge in the flow
  - Can find a path faster than existing algs for graphs with only short paths
  - Can figuring out how to sabotage connections in graphs with high throughput bridges
  - Other applications?

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- Recent work (with Stacey Jeffery and Alvaro Piedrafita): algorithm to sample an edge in the flow
  - Can find a path faster than existing algs for graphs with only short paths
  - Can figuring out how to sabotage connections in graphs with high throughput bridges
  - $\circ$  Other applications?
- When is path detection reduction approach optimal?
- Way to find optimal weights?

#### Thank you! **Funding: Collaborators:**





Da-Yeon Koh



Chloe Ye



Teal Kai Witter DeLorenzo



**Stacey Jeffery** 

Anderson





Michael Jarret



Jay-U Chung



Alvaro Piedrafita