# Problems with Multiple Oracles

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#### What '



IT PROMISES TO SOLVE SOME OF HUMANITY'S MOST COMPLEX PROBLEMS. IT'S BACKED BY JEFF BEZOS, NASA AND THE CIA. EACH ONE COSTS \$10,000,000 AND OPERATES AT 459° BELOW ZERO. AND NOBODY KNOWS

HOW IT ACTUALLY WORKS

THE INFINITY MACHINE



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### What to do with a Quantum Computer?

• Let's try to solve something hard: 3-SAT

Traveling Salesman Problem [Karp '72] Tetris [Demaine et al '03]

Is a graph planar? [Grigoriev et al '07] Scheduling jobs [Ullman '75]

> Sudoku [Yato et al ]

Graph Coloring [Karp '72]

# **3-SAT**

• Goal: Want to satisfy a set of Boolean clauses, each with 3 variables.

$$(x_1 \vee \neg x_2 \vee x_3) \wedge (x_2 \vee x_3 \vee \neg x_4)$$

Each variable  $x_i$  can take value 0 or 1

V is logical OR:  

$$0 \land 0 = 0$$
  
 $0 \land 1 = 1$   
 $1 \land 0 = 1$   
 $1 \land 1 = 1$   
 $\neg$  is logical NOT:  
 $\neg 0 = 1$   
 $\neg 0 = 1$   
 $(\land 0 = 0)$   
 $1 \land 0 = 0$   
 $1 \land 0 = 0$   
 $1 \land 0 = 0$   
 $1 \land 1 = 1$ 

-1 = 0

### **3-SAT**

Is there an assignment of the *n* variables  $\{x_1, x_2, \dots, x_n\}$  such that  $F(x_1, x_2, \dots, x_n) = 1$ , where

$$F(x_1, x_2, \cdots, x_n) = (x_1 \lor \neg x_3 \lor x_7) \land (x_2 \lor x_3 \lor \neg x_6) \land (x_1 \lor x_7 \lor x_{10}) ..$$
  
~poly(n) clauses (e.g.  $Cn^2$ )

#### **Algorithm for 3-SAT**

$$F(x_1, x_2, \cdots, x_n) = (x_1 \vee \neg x_3 \vee x_7) \wedge (x_2 \vee x_3 \vee \neg x_6) \wedge (x_1 \vee x_7 \vee x_{10}) \dots$$
  
~poly(n) clauses (e.g.  $Cn^2$ )

Guess a satisfying assignment. Test if all clauses are satisfied
 Need to test ~2<sup>n</sup> possible inputs. With quantum computer can do in  $\sqrt{2^n}$  steps

# Outline

- Oracles and Oracle Models
- Related work
- Simple Example: Search with Multiple Oracles
- Open Problems and Directions for Future Work

### **Standard Oracle Model vs 3-SAT**

3-SAT	Oracle Model
Given a function $F(x_1, x_2, \dots, x_n) = (x_1 \lor \neg x_3 \lor x_7)$	Initially function $F$ is unknown
Determine if input $x$ exists such that $F(x) = 1$	<ul> <li>Determine some property of <i>F</i></li> <li>Does <i>x</i> exists such that <i>F</i>(<i>x</i>) = 1?</li> <li>Is <i>F</i> one-to-one?</li> </ul>

### **Standard Oracle Model**

Goal: Determine a property of a function F(x) for Boolean input  $x = \{x_1, x_2, \dots, x_n\}$ , given an oracle for F.



Only care about # of oracle calls (queries)

# **Multiple Oracles with Costs Model**

Goal: Determine a property of a function F(x) for Boolean input  $x = \{x_1, x_2, \dots, x_n\}$ , given a set of oracles associated with functions  $\{F_1, \dots, F_k\}$  which each have some information related to F

Care about total cost =  $\sum_{i=1}^{k} q_i c_i$  where  $q_i$  is the # of times Oracle *i* is used

# **Related Work**

- Ambainis '10: One oracle, but querying different *x* requires different amounts of time
  - E.g. To learn  $f(00 \cdots 00)$  takes time 1, but to learn  $f(11 \cdots 11)$  takes time 2
- Montanaro '09: Searching when given some additional information as to the location of the marked item.
  - E.g. Told that  $f(00 \cdots 00) = 1$  is more likely than  $f(11 \cdots 11)$
- Cerf et al. '00: Use multiple oracles to speed up evaluation of satisfiability problems.
  - No cost, No lower bounds, Need certain structure.

# Searching with an Oracle

Goal: Determine x such that F(x) = 1. Can ask oracle, "Is the *i*<sup>th</sup> item the starred item?"

- Classically:  $\sim N$  queries to oracle
- Quantumly:  $\sim \sqrt{N}$  queries to oracle [Grover '97, Bennett et al. '97, Zalka '99]









Can ask  $\bigstar$  Oracle, "Is the *i*<sup>th</sup> item starred?" with cost  $c_{\star}$ Can ask Oracle, "Is the *i*<sup>th</sup> item striped?" with cost  $c_{\parallel}$ 

**C**★>**C**∣

Promised: The starred item is also striped







Can ask  $\bigstar$  Oracle, "Is the *i*<sup>th</sup> item starred?" with cost  $c_{\star}$ Can ask () Oracle, "Is the *i*<sup>th</sup> item striped?" with cost  $c_{\parallel}$ 

C<sub>\*</sub>>C|

Promised: The starred item is also striped



Sometimes best to check all N items using 4 Oracle

Otherwise can check all *N* items using cheaper oracle, but still need to use  $\bigstar$  Oracle for *M* items

### **Quantum Query Complexity Bounds**



Size of Problem

#### • TOOL:

Need at least  $\sim \sqrt{N}$  queries to quantum oracle to find one item out of N [Bennett et al '97]

\* "at least  $\sim \sqrt{N}$ " means at least  $A\sqrt{N}$  for some constant A as N gets large. =  $\Omega(\sqrt{N})$ 

• How much can **Oracle** help?



• How much can **Oracle** help?



- How much can **Oracle** help?
- At best, can narrow down to *M* items. But now searching for 1 among *M*, need at least  $\sim \sqrt{M}$  queries to **Cracle TOOL**



• How much can **Oracle** help?

At best, can narrow down to *M* items. But now searching for 1 among *M*, need at least  $\sim \sqrt{M}$  queries to **Cracle TOOL** 



★ Oracle has cost  $c_{\star}$ , and need to use it ~ $\sqrt{M}$  times. Always will have a cost of ~ $c_{\star}\sqrt{M}$ 

- Need  $\bigstar$  Oracle  $\sim \sqrt{M}$  times; How many times do we need  $\blacksquare$  Oracle?
- What if only needed to use ( ) **Oracle**  $\sqrt{M}$  times? (For contradiction.)
  - This would be awesome!
  - Would always want to have an **()** Oracle, because it helps so much

Idea: Even if aren't given **()** Oracle, create it using + Oracle.











• What if only needed to use ( ) Oracle  $\sqrt{M}$  times? (For contradiction.) Idea: Even if aren't given ( ) Oracle, create it using  $\bigstar$  Oracle.

Can simulate **Oracle** using one (two) queries to **Cracle**!



Only need to use  $\bigstar$  Oracle  $\sim \sqrt{M}$  times to simulate  $\blacksquare$  Oracle.

• What if only needed to use ( ) Oracle  $\sqrt{M}$  times? (For contradiction.) Idea: Even if aren't given ( ) Oracle, create it using  $\bigstar$  Oracle.

Can simulate  $\bigcirc$  **Oracle** using one (two) queries to  $\bigstar$  **Oracle**!

Only need to use  $\bigstar$  Oracle  $\sim \sqrt{M}$  times to simulate  $\blacksquare$  Oracle.



Previously showed need  $\sim \sqrt{M}$  queries to  $\Rightarrow$  Oracle if have an  $\bigcirc$  Oracle.

• What if only needed to use ( ) Oracle  $\sqrt{M}$  times? (For contradiction.) Idea: Even if aren't given ( ) Oracle, create it using  $\bigstar$  Oracle.

Can simulate  $\bigcirc$  **Oracle** using one (two) queries to  $\bigstar$  **Oracle**!

- Only need to use  $\bigstar$  Oracle  $\sim \sqrt{M}$  times to simulate  $\blacksquare$  Oracle.
- Previously showed need  $\sim \sqrt{M}$  queries to  $\neq$  Oracle if have an  $\bigcirc$  Oracle.

Can find starred item using  $\sim \sqrt{M}$  queries to  $\bigstar$  Oracle

• What if only needed to use ( ) Oracle  $\sqrt{M}$  times? (For contradiction.) Idea: Even if aren't given ( ) Oracle, create it using  $\bigstar$  Oracle.

Can simulate **Oracle** using one (two) queries to **Cracle** 

- Only need to use  $\bigstar$  Oracle  $\sim \sqrt{M}$  times to simulate  $\blacksquare$  Oracle.
- Previously showed need  $\sim \sqrt{M}$  queries to  $\Rightarrow$  Oracle if have an  $\blacksquare$  Oracle.

M queries to 🗡 🛛

• Using this argument:



Always need to use either  $\bigstar$  Oracle or () Oracle at least  $\sim \sqrt{N}$  times.



# Algorithm for Searching with Multiple Oracles

Amplitude Amplification

# **Quantum Query Complexity Bounds**



Size of Problem

### **Algorithm for 3-SAT**

$$F(x_1, x_2, \cdots, x_n) = (x_1 \lor \neg x_3 \lor x_7) \land (x_2 \lor x_3 \lor \neg x_6) \land (x_1 \lor x_7 \lor x_{10}) \dots$$
  
~poly(n) clauses (e.g.  $Cn^2$ )

- Guess a satisfying assignment. Test if all clauses are satisfied EXPENSIVE
   Need to test ~2<sup>n</sup> possible inputs. With quantum computer can do in √2<sup>n</sup> steps
- Guess a satisfying assignment. Test if ~log(n) clauses are satisfied CHEAP
   Defines a subset of possible solutions, including the true satisfying assignment, if it exists
   What is M?

# **Directions for Future Work**

Create tight bounds for searching with multiple oracles

 Adversary Bound/Span programs
 Geometric picture

- Can we create a general framework for understanding oracles with costs, in the way that the adversary bound is a framework for understanding standard oracle problems
- Many quantum oracle problems does it make sense to add multiple oracles to these problems?

# **Classical Algorithm**

- 1. Choose item at random and test if striped using Oracle 2
- 2. If it is striped, test if starred using Oracle 1

Worst case cost:

 $c_1(M-1) + c_2N$