

Characterizing Quantum Operations

Shelby Kimmel

Joint work with Marcus Silva,
Raytheon BBN Technologies

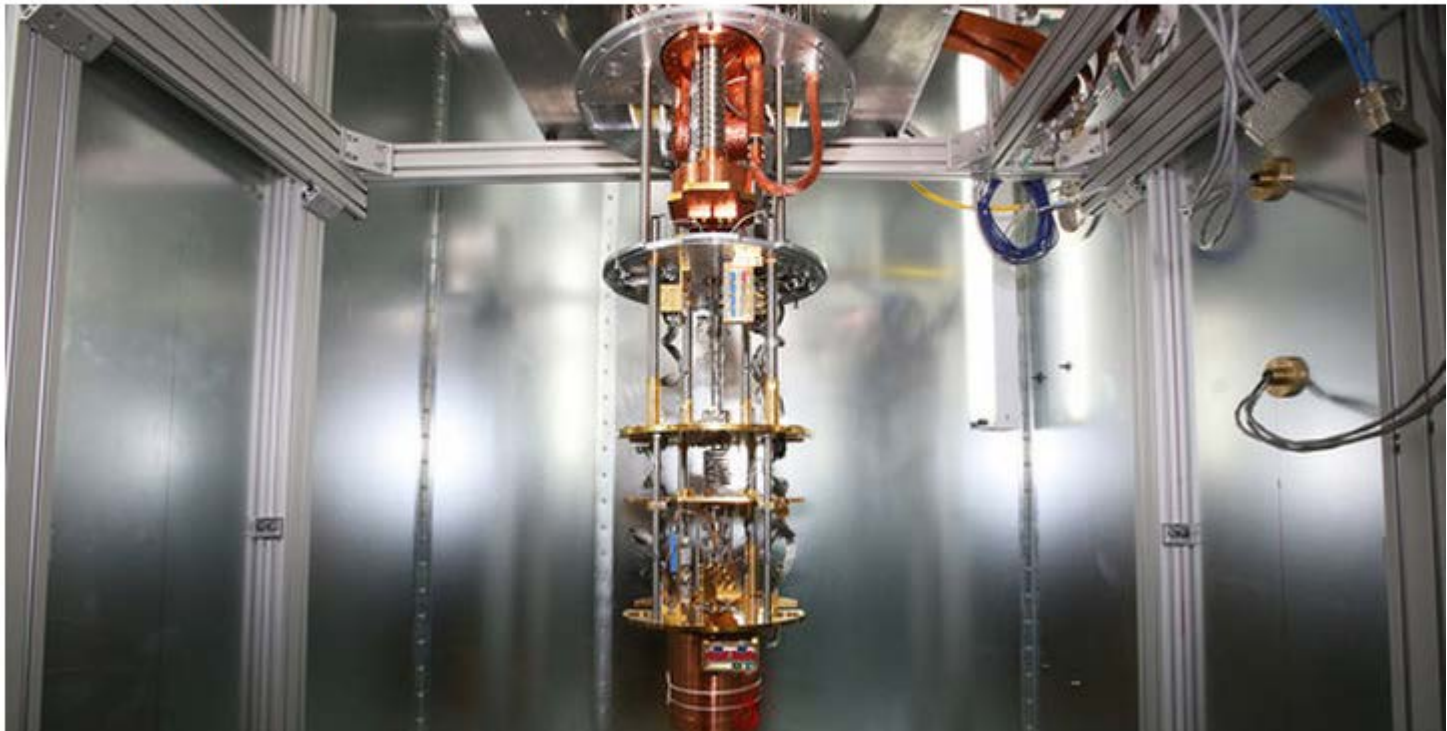
Quantum Computers Will Be Cool

The New York Times

Business Day

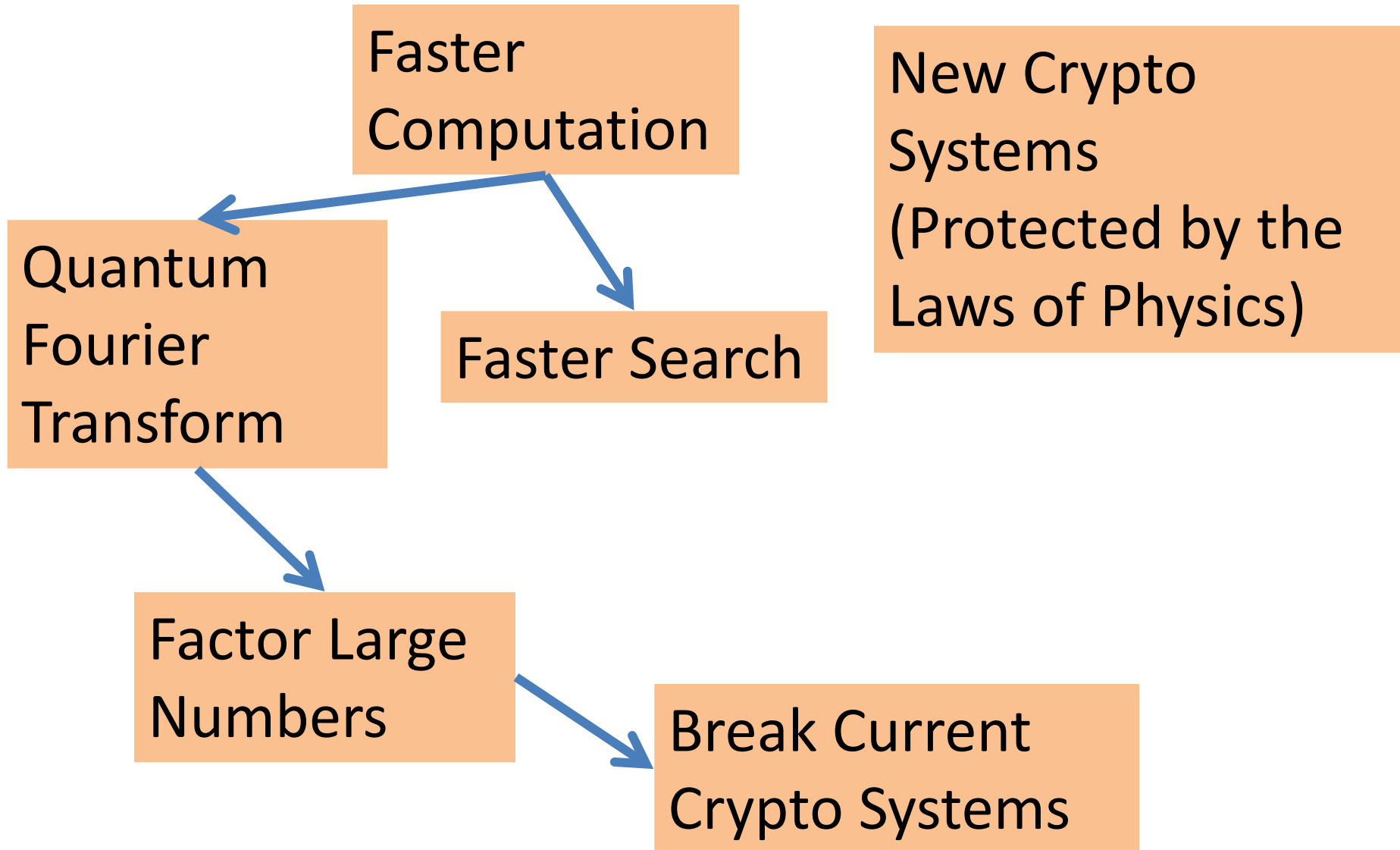
Technology

A Strange Computer Promises Great Speed



By [QUENTIN HARDY](#), Published: March 21, 2013

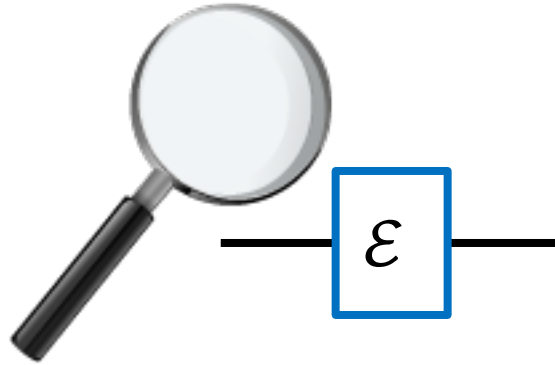
Quantum Computers Will Be Cool



Why Don't We All Have One?

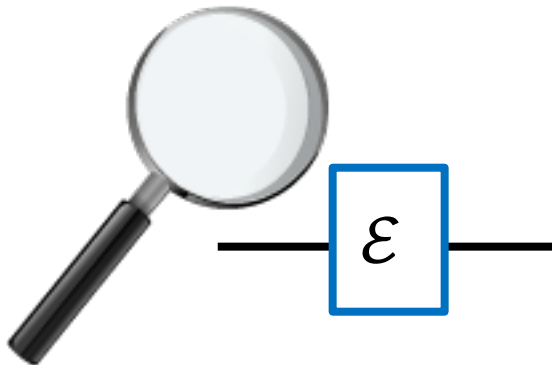
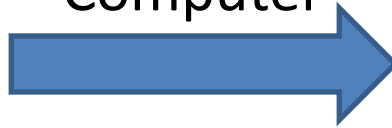
Too Many Errors

Can Improve Operations with Better Characterization of Errors



“Depolarizing error”

Improvement to
Computer

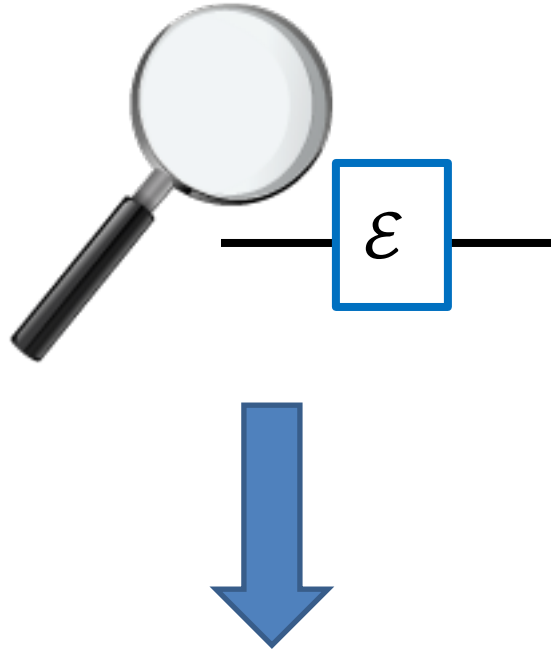


“Extra rotation around z-axis”

Improvement to
Computer



Standard Techniques Have Problems



Often, outcome is not a
valid quantum operation!

Outline: Improving Quantum Process Characterization

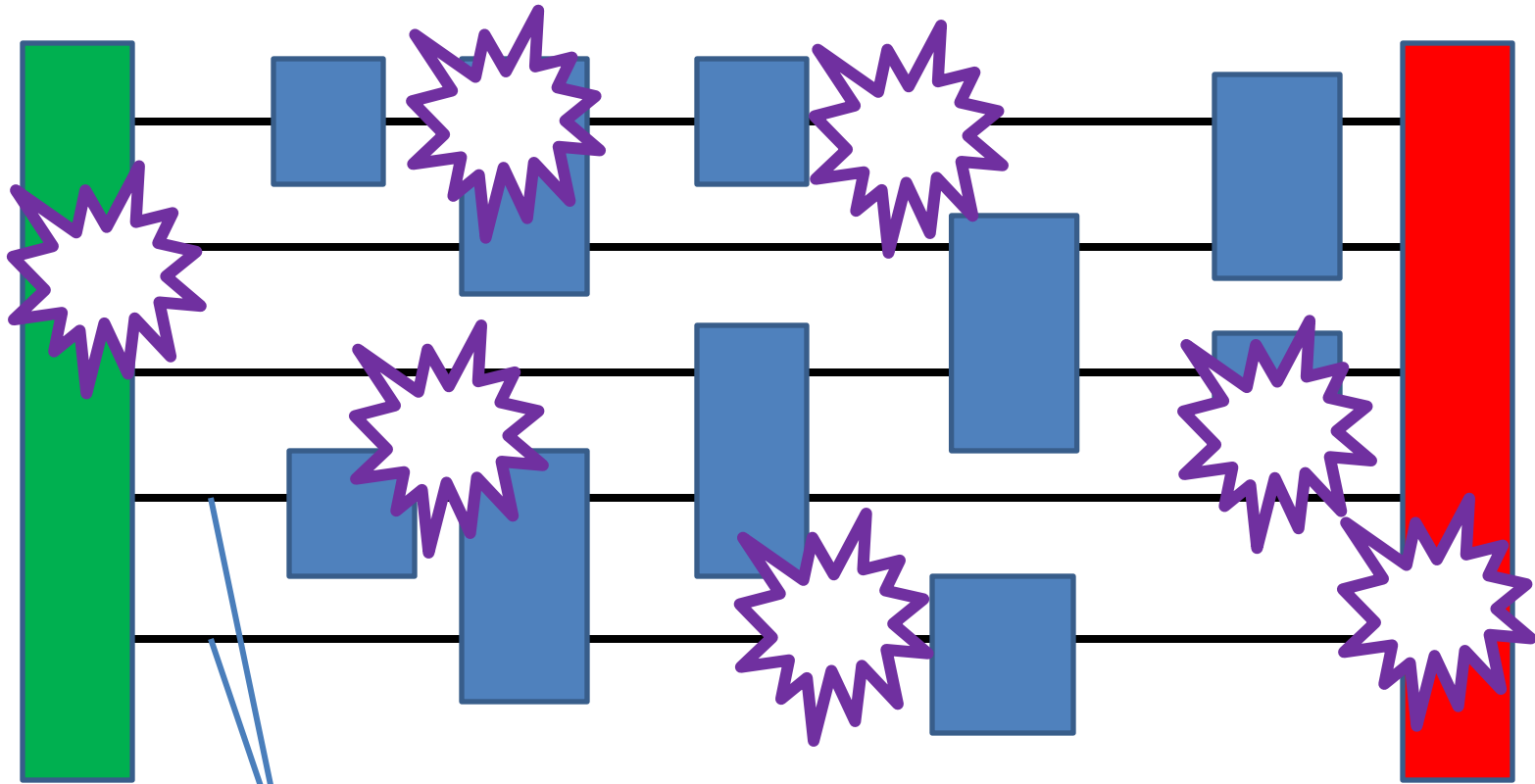
- **Motivation:** Process Characterization can improve errors and thus lead towards the realization of a quantum computer
- **Problem:** Standard techniques for process characterization have systematic errors
- **Solution:** We provide a way to get around these systematic errors and get lots of useful information about quantum operations.
(Magesan et al. '11, '12, Kimmel and Silva...soon!)

Quantum Computation (Circuit Model)

State Preparation

Operations

Measurement

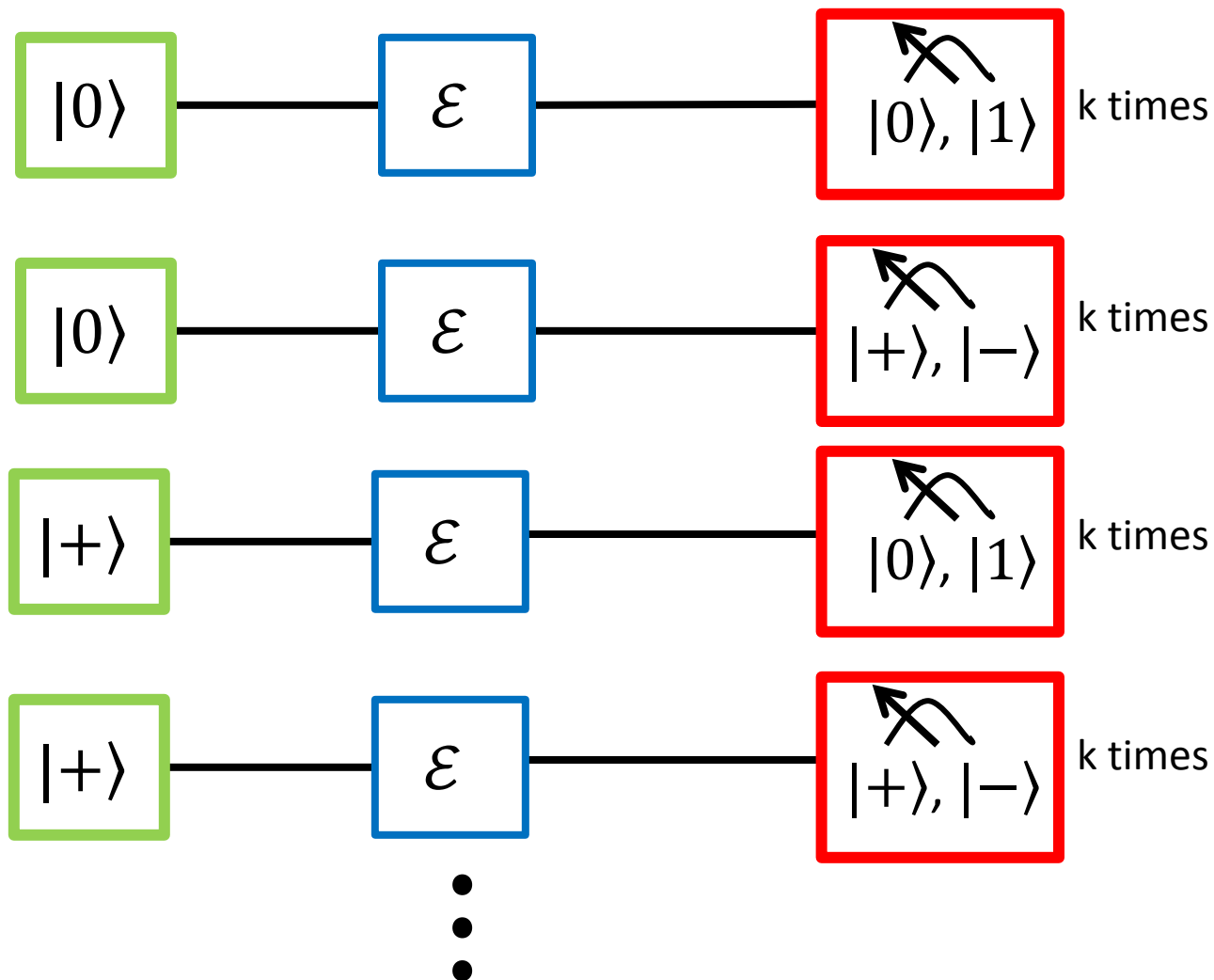


Qubits

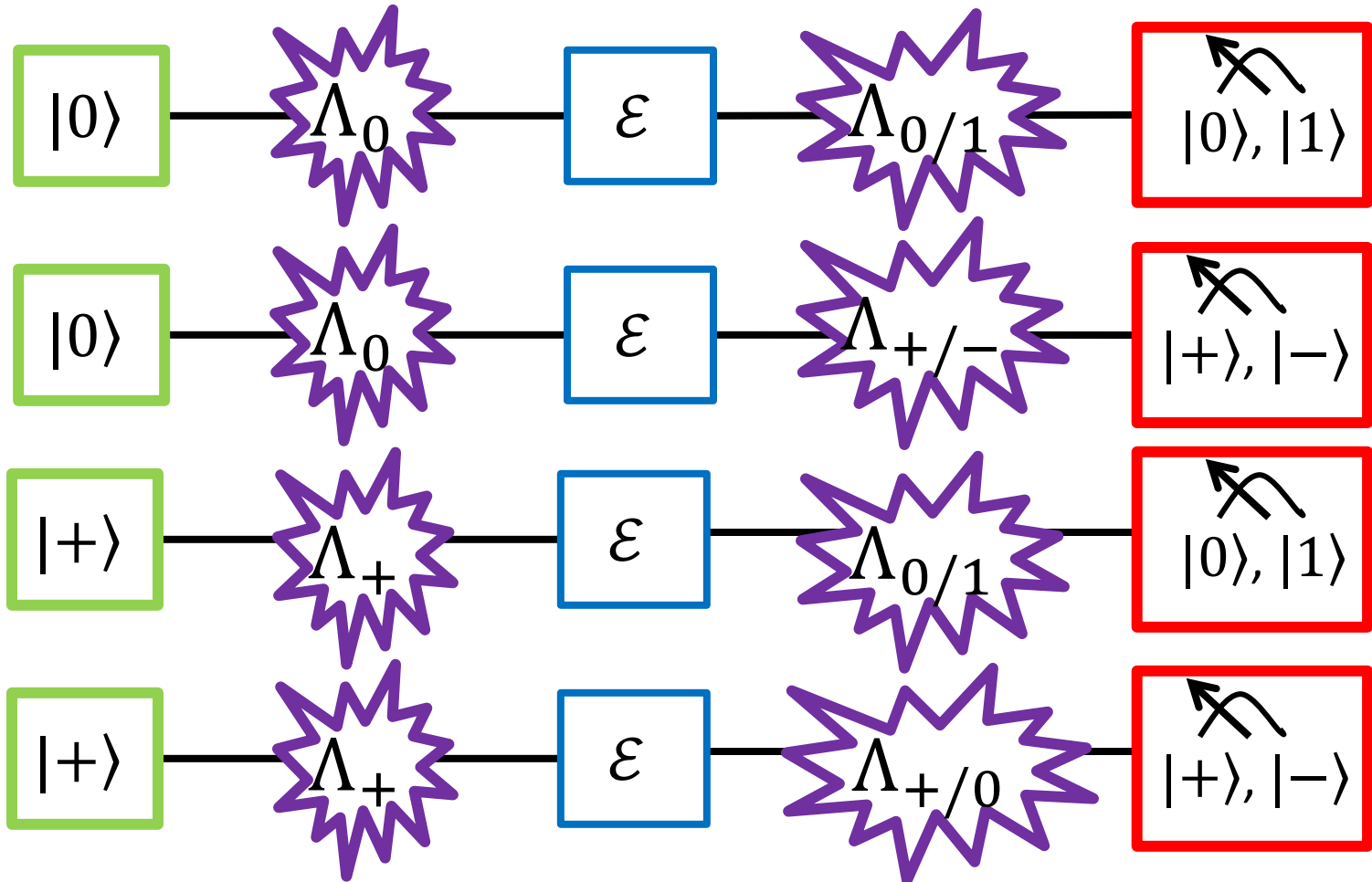
Quantum Process

- Completely positive trace preserving (CPTP) map = any process that takes valid quantum states to valid quantum states.
- E.g. unitary, depolarizing process, dephasing process, amplitude damping process
- n qubits, $16^n - 4^n$ free parameters

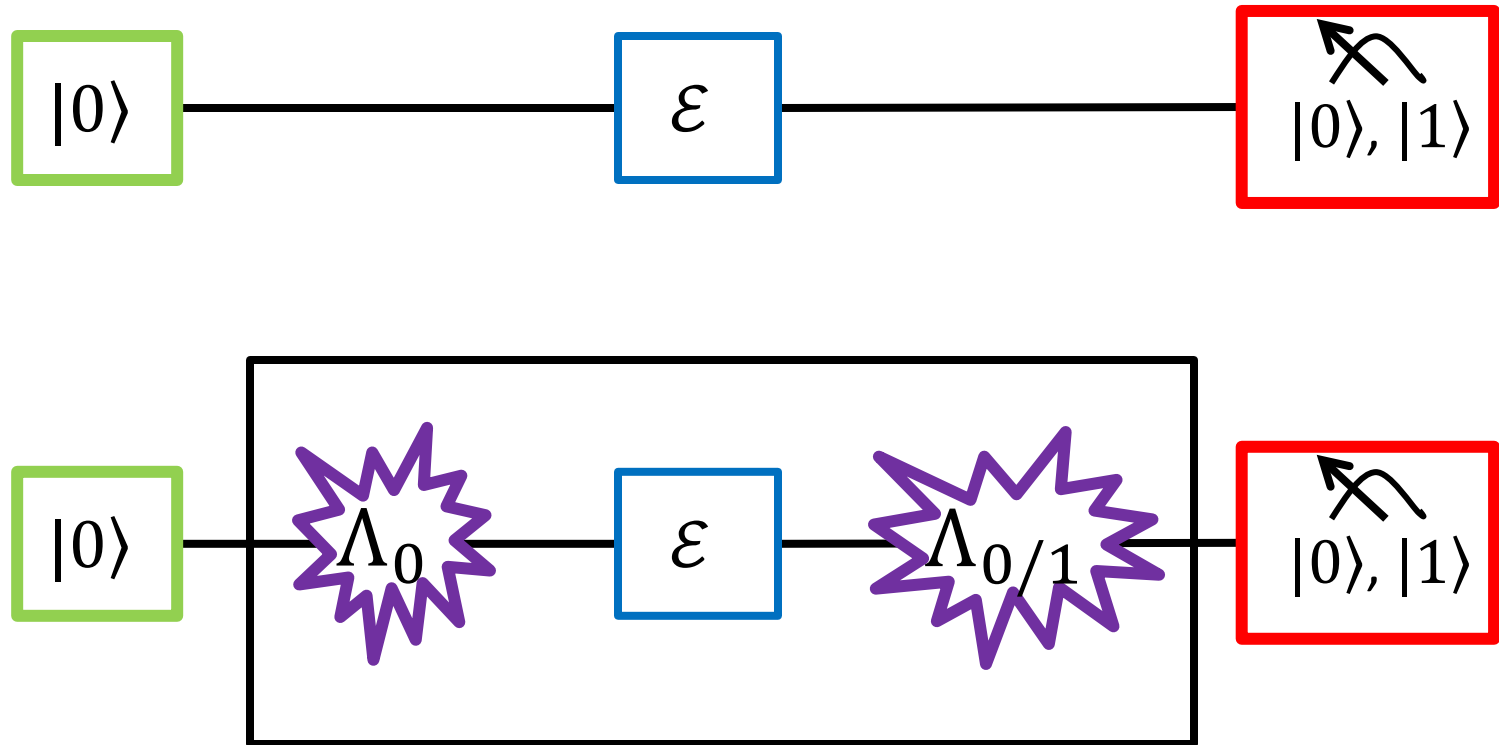
Standard Quantum Process Characterization



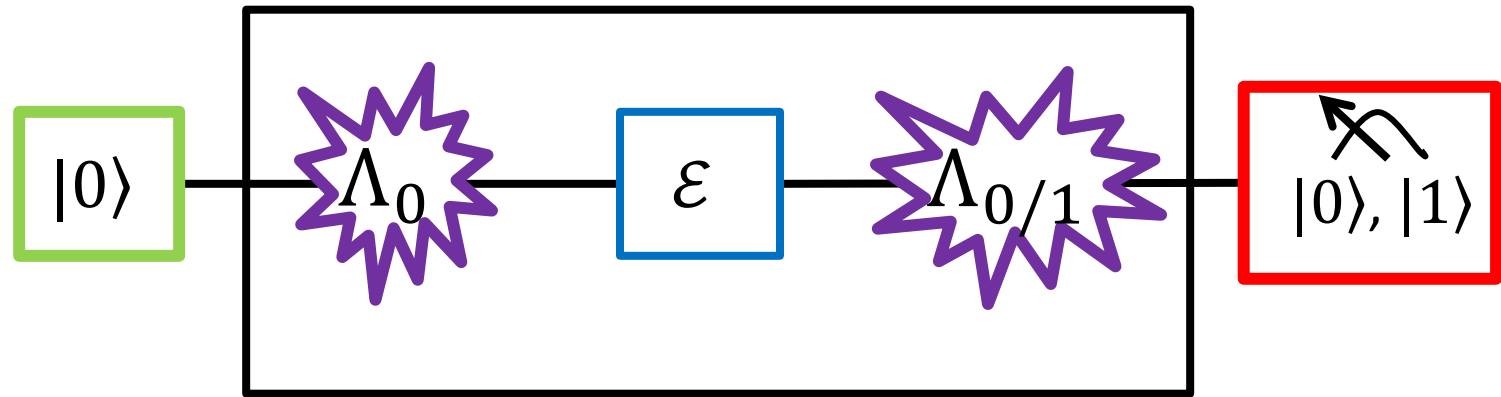
The Problem



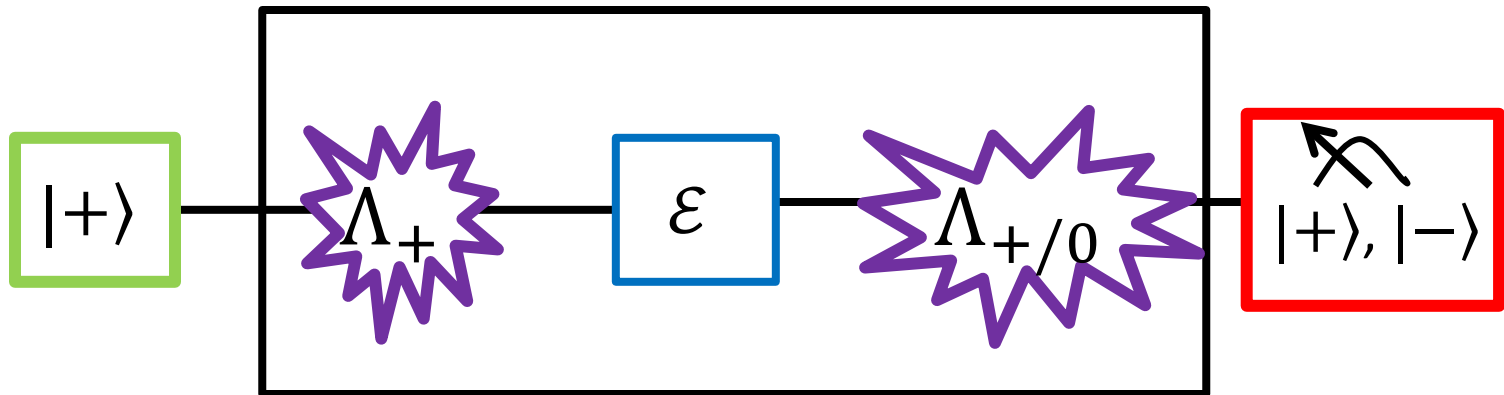
The Problem



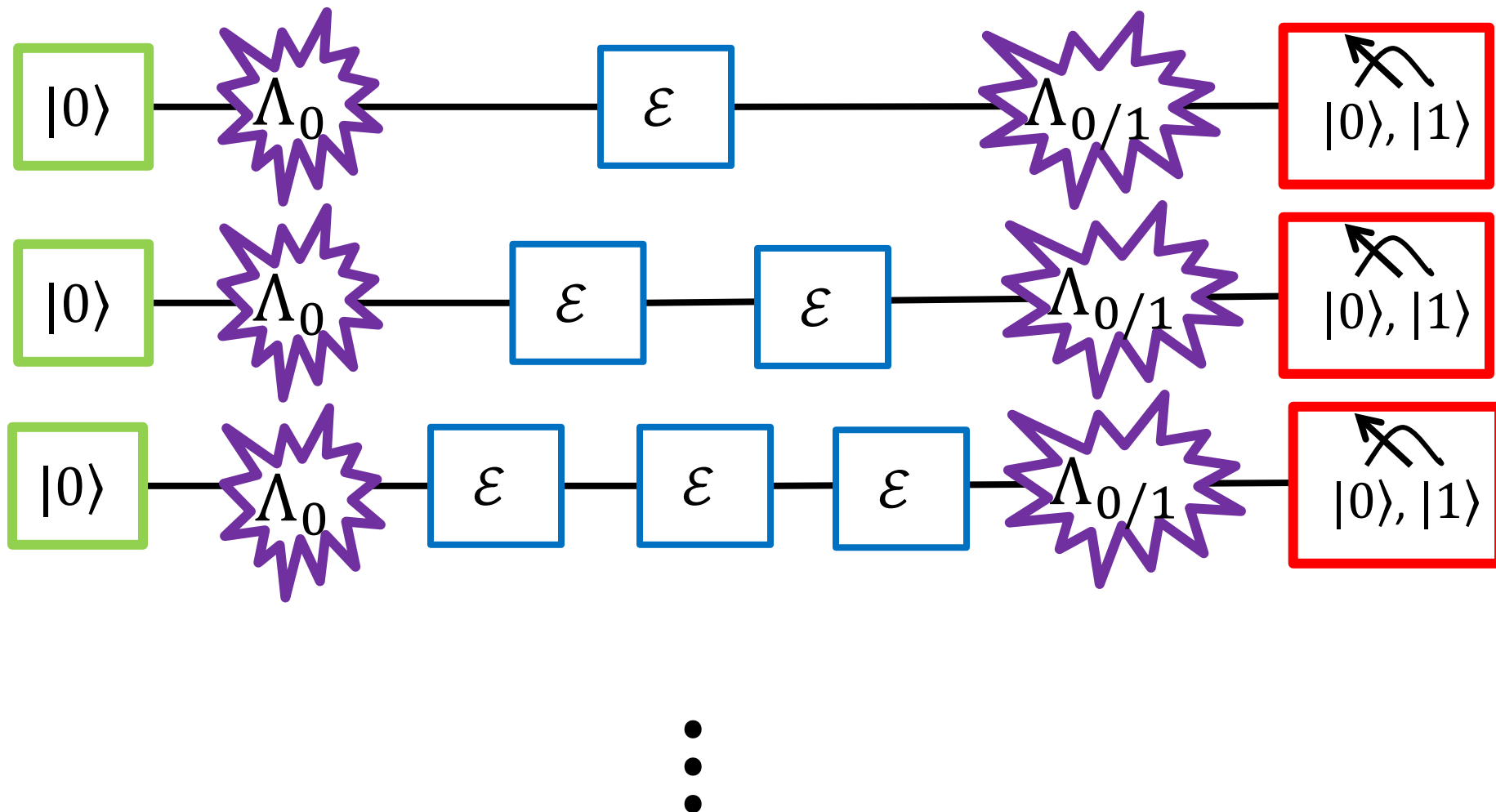
The Problem



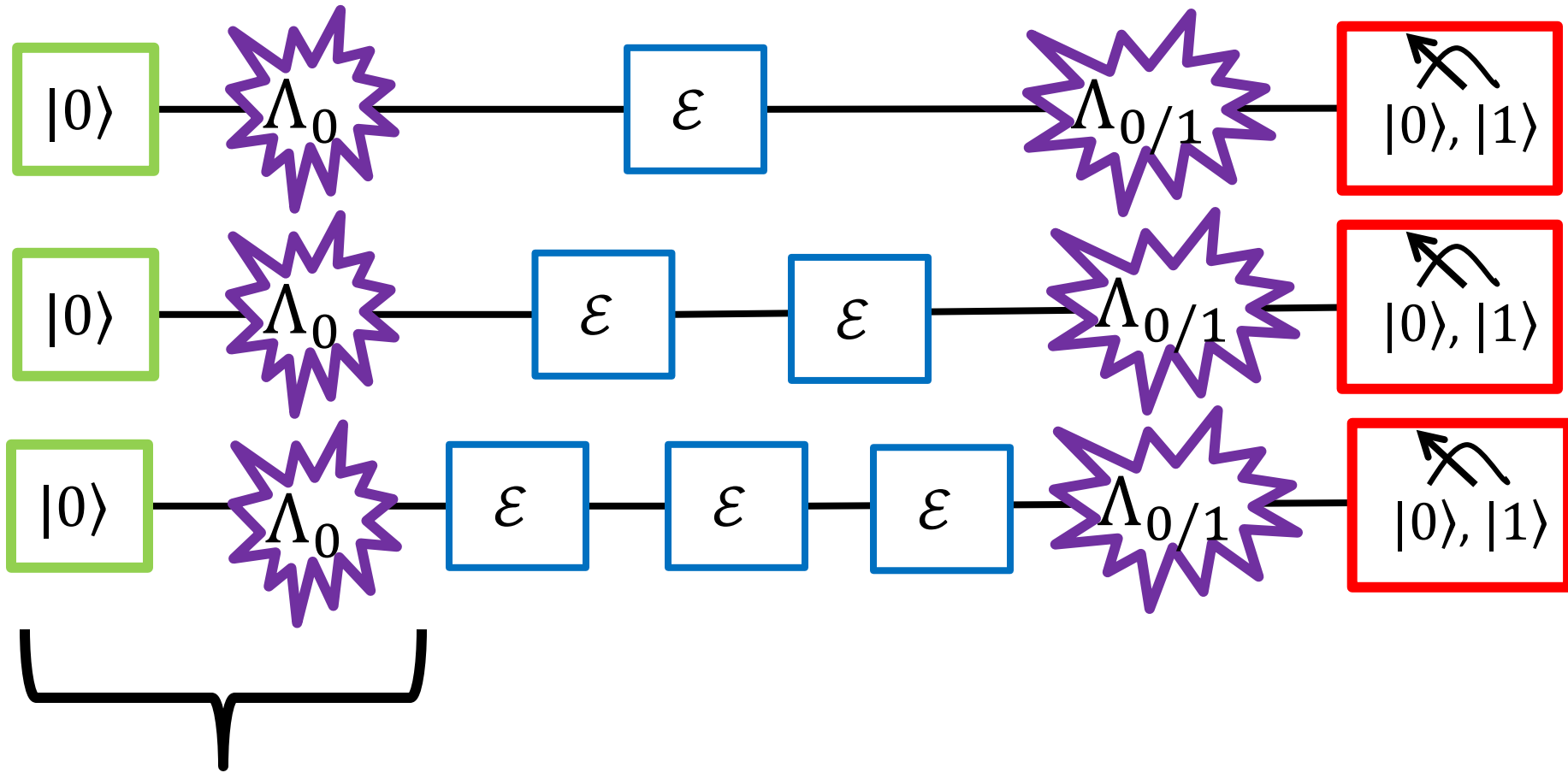
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Repeated Application

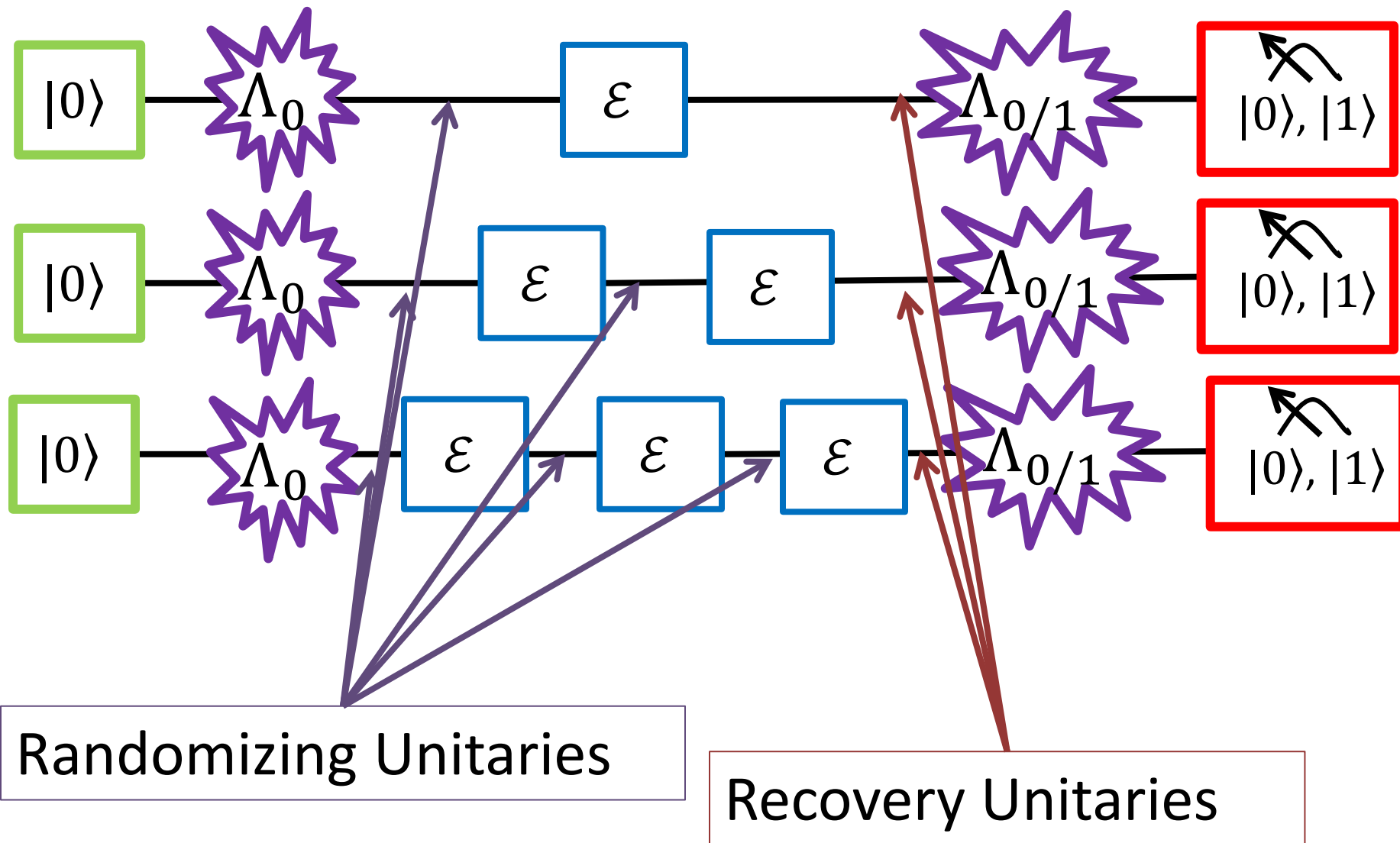


Repeated Application

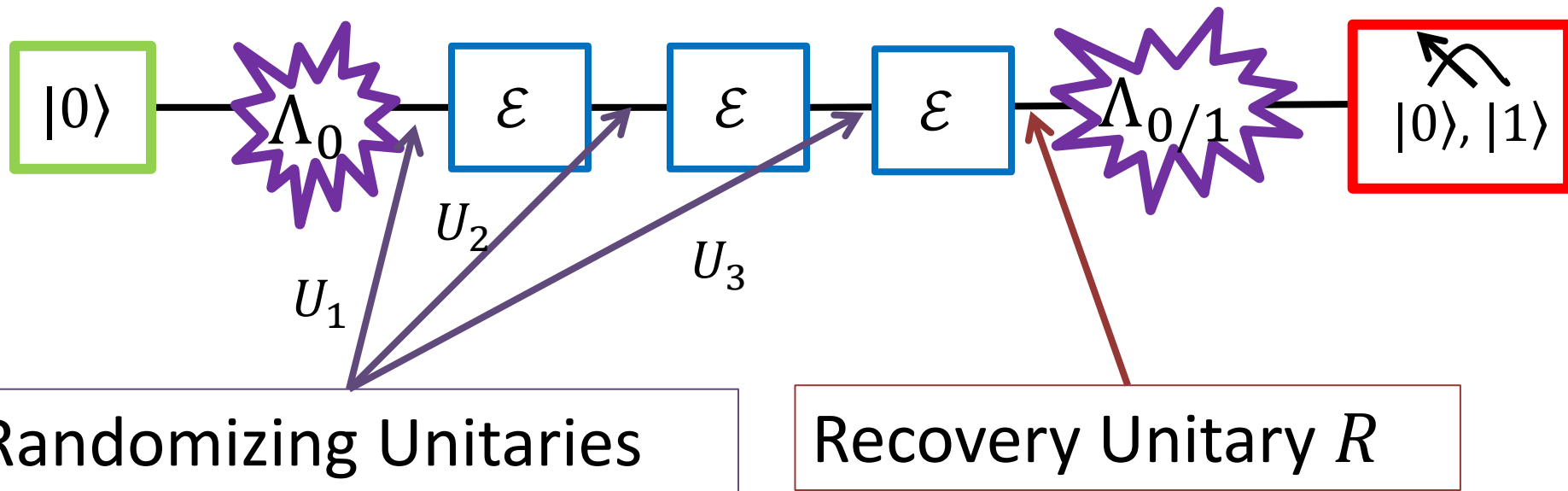


If eigenstate of \mathcal{E} , will only see how \mathcal{E} acts on *this* state

Repeated Application



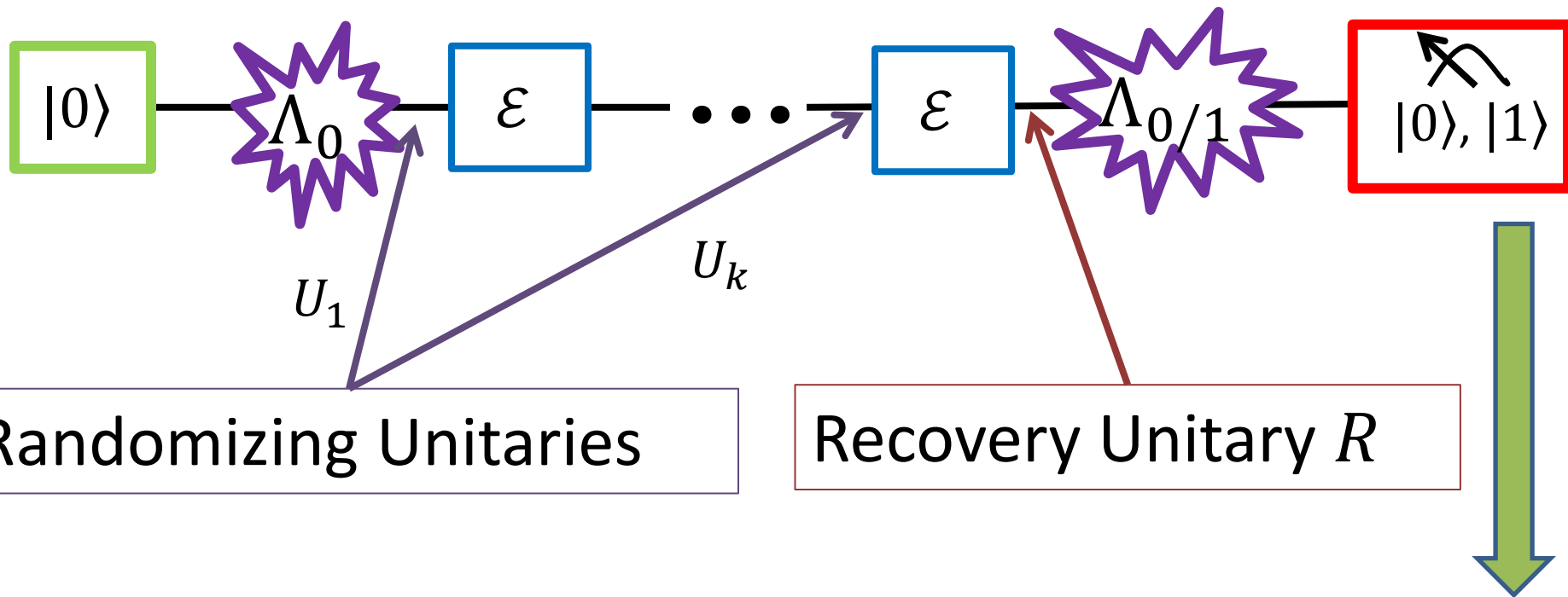
Recovery Operation



Pick a unitary V to compare \mathcal{E} with. Then

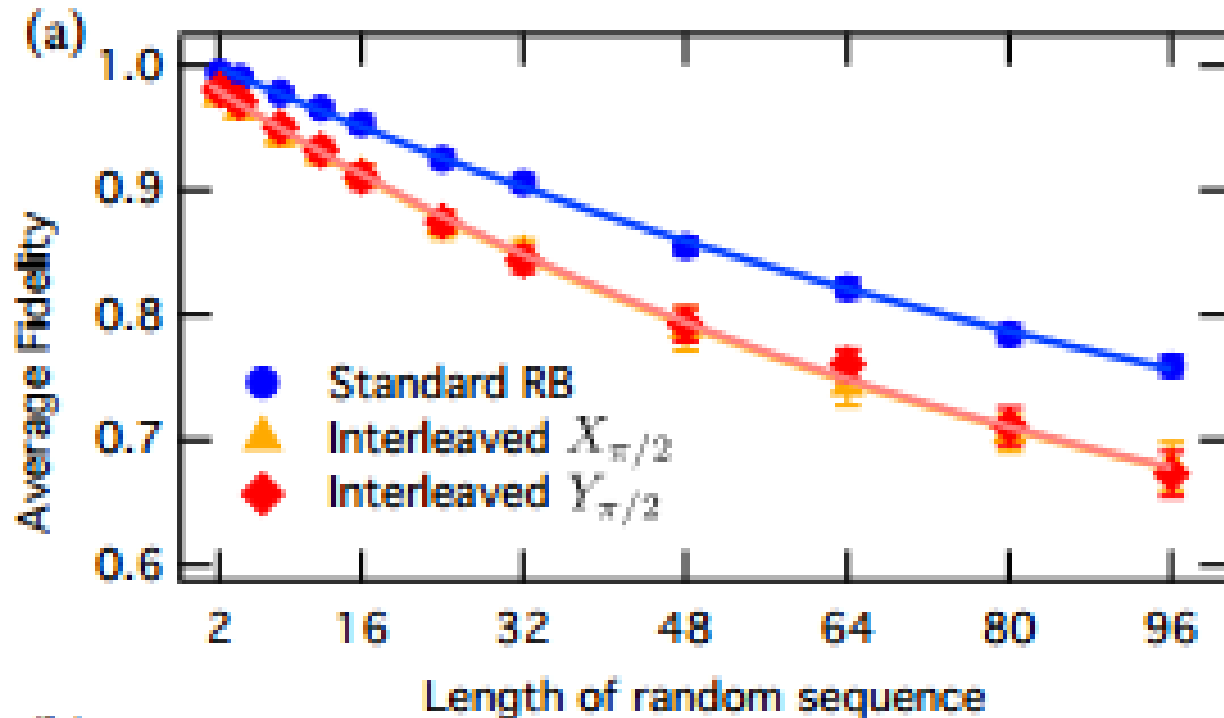
$$R = (U_1 V U_2 V U_3 V)^\dagger$$

Randomizing Operations



Prob. of outcome 0: $Ap^k + B$
 A, B depend on $\Lambda_0, \Lambda_{0/1}$
 p depends on overlap of \mathcal{E} with V

Randomizing Operations

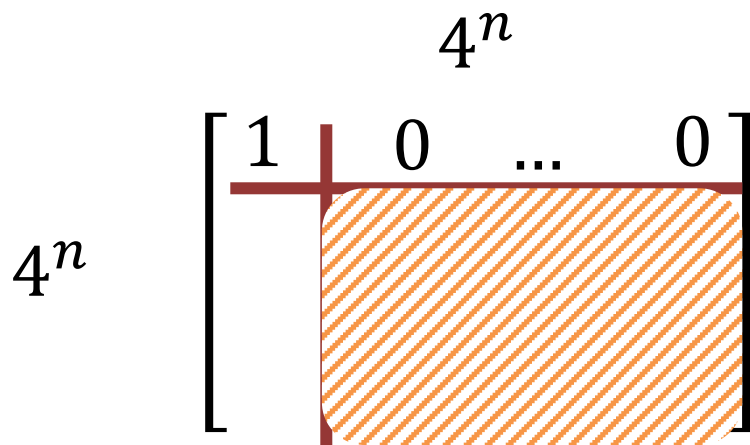


Magesan et al. '12

Prob. of outcome 0: $Ap^k + B$
 A, B depend on $\Lambda_0, \Lambda_{0/1}$
 p depends on overlap of \mathcal{E} with V

“Overlap of \mathcal{E} with V ?”

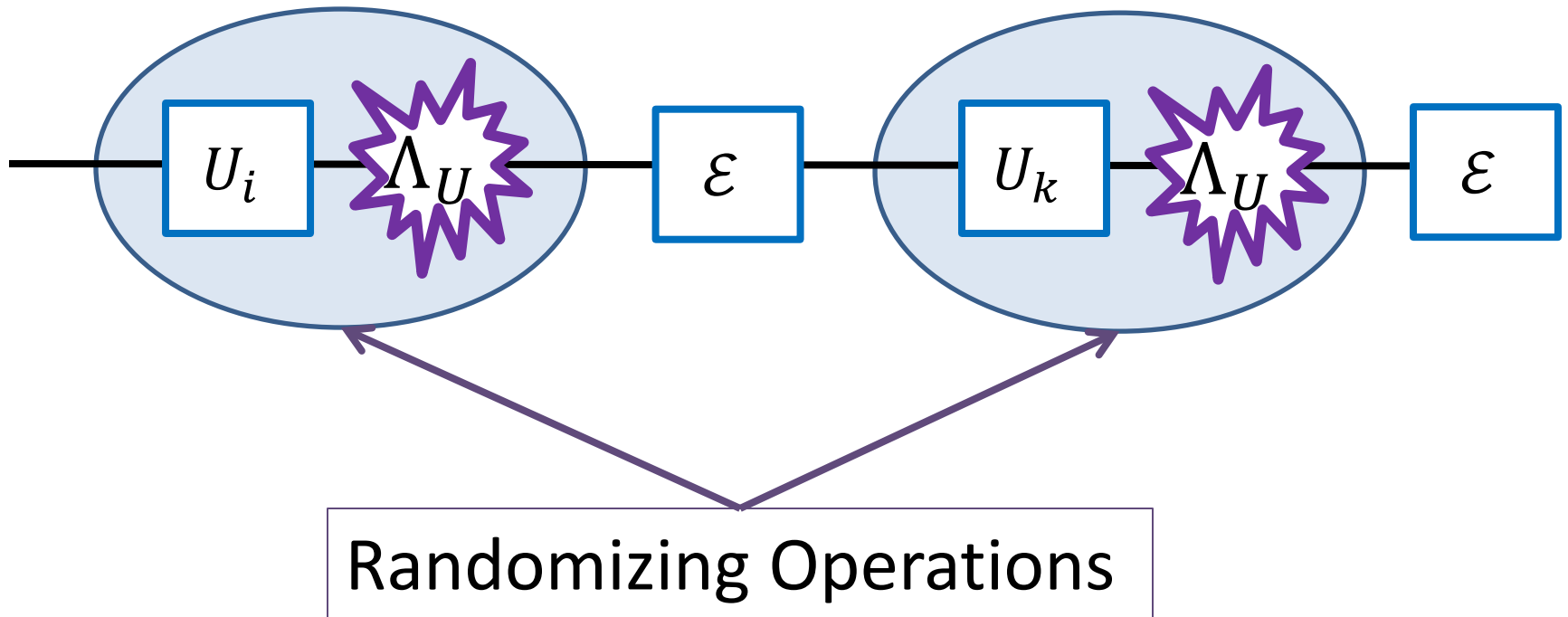
CPTP map: $16^n - 4^n$ parameters for n qubit map



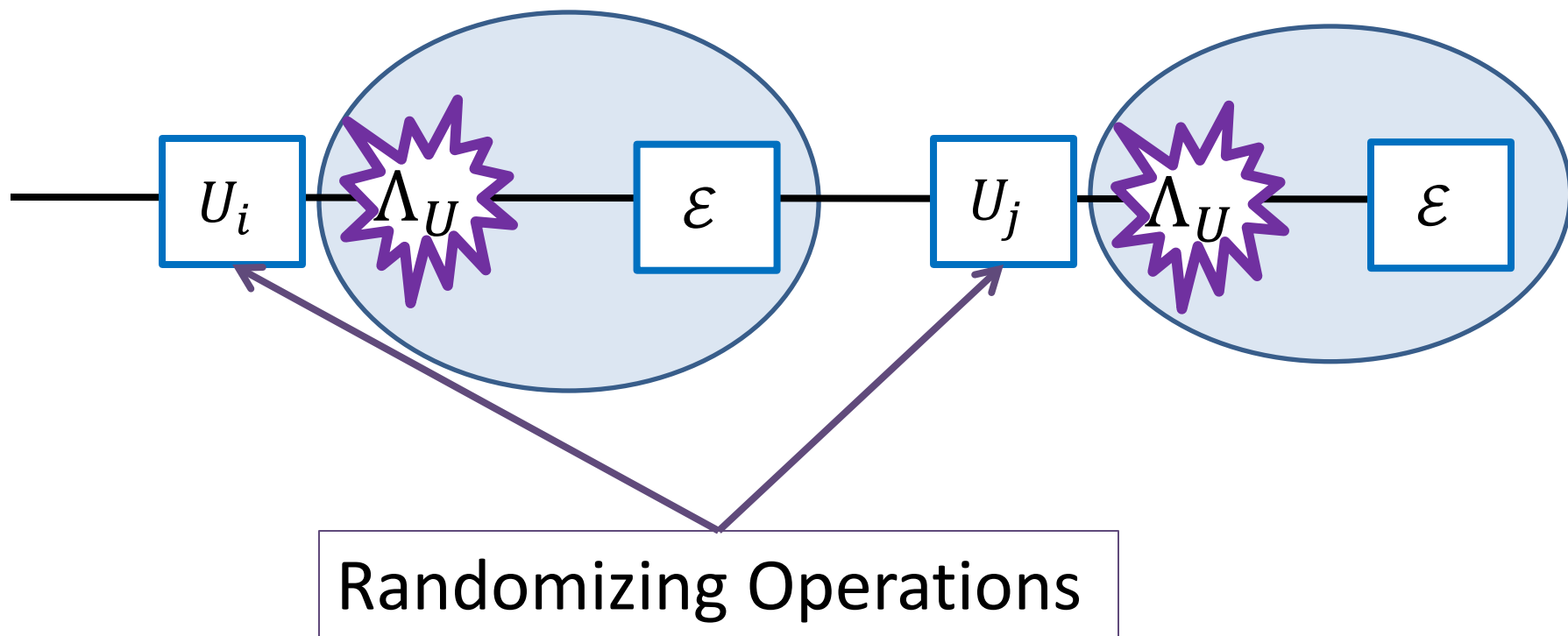
Choose many different V 's and measure overlaps. Each overlap gives a new parameter of this matrix!

We learn: $16^n - 2 * 4^n + 1$ parameters

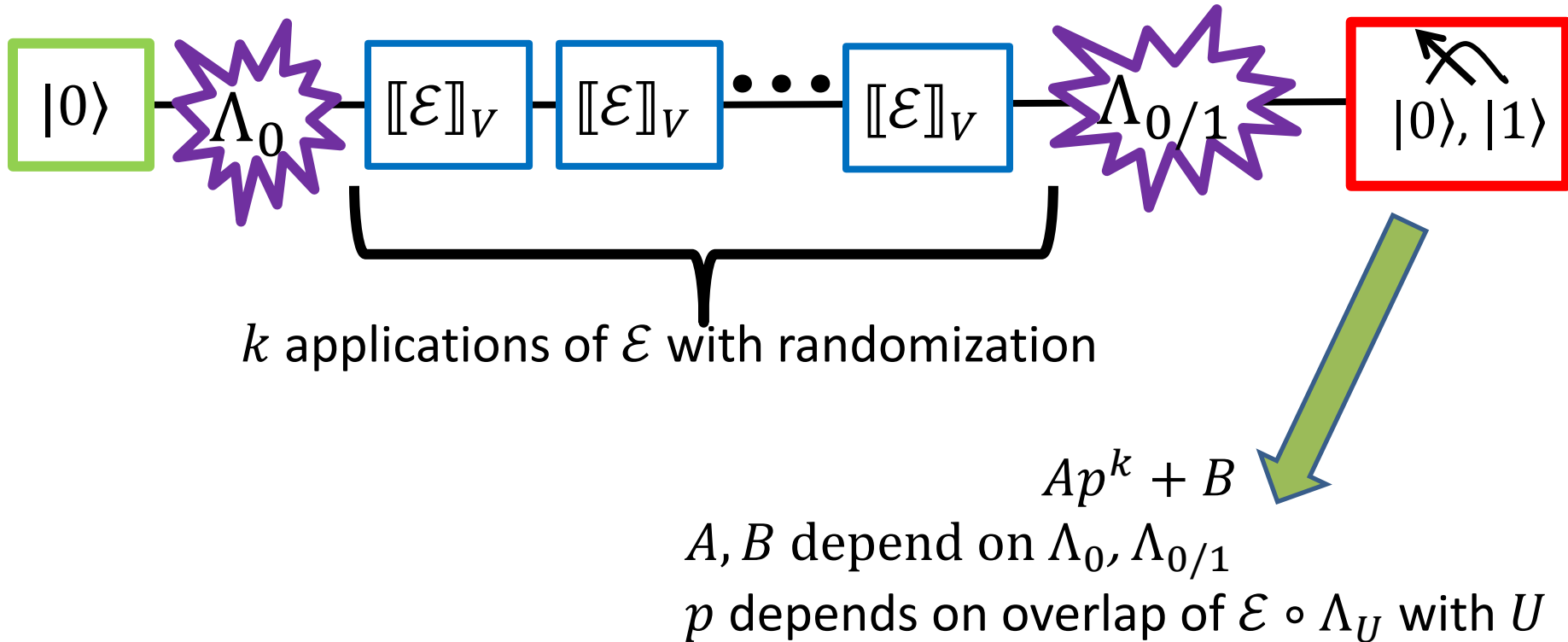
Errors!



Errors!



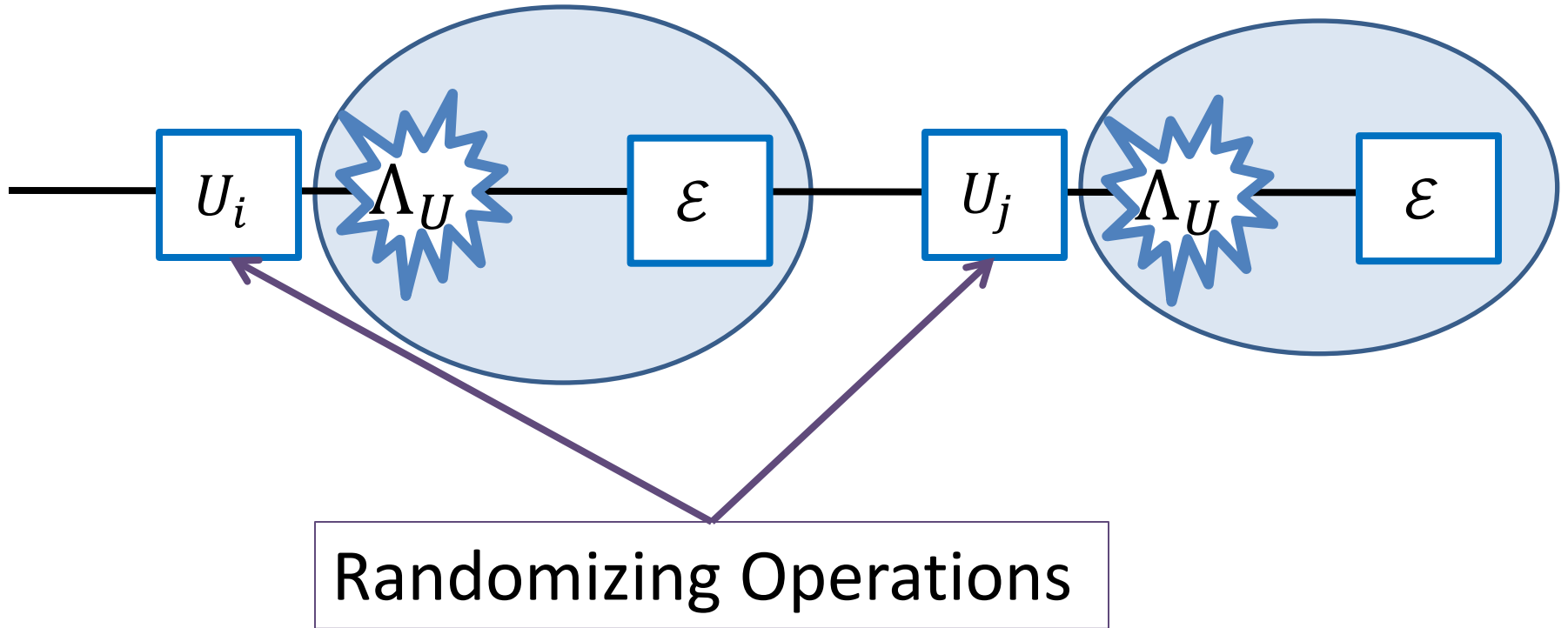
Randomizing Operations with Error



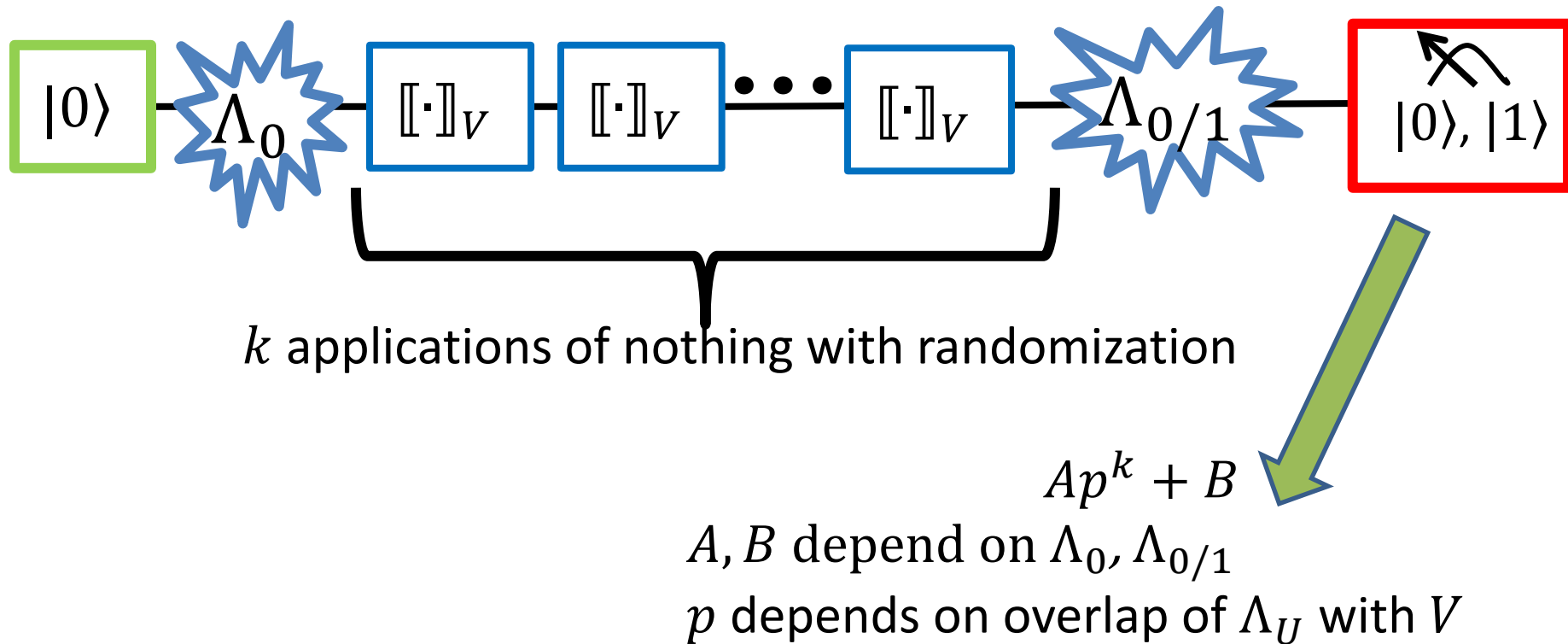
Trading One Error For Another?

- Originally had state preparation and measurement errors complicating things
- Now have randomizing errors complicating things
- BUT – we can characterize randomizing errors!

Errors!



Randomizing Operations with Error



Main Result

$$\begin{array}{l} \text{almost complete characterization of } \Lambda_U \\ + \\ \text{almost complete characterization of } \mathcal{E} \circ \Lambda_U \\ = \\ \text{almost complete characterization of } \mathcal{E} \end{array} \quad \begin{array}{l} \left[\begin{array}{c|ccc} 1 & 0 & \dots & 0 \\ \hline & \text{orange diagonal} & & \end{array} \right] \\ \\ \left[\begin{array}{c|ccc} 1 & 0 & \dots & 0 \\ \hline & \text{green diagonal} & & \end{array} \right] \\ \\ \left[\begin{array}{c|ccc} 1 & 0 & \dots & 0 \\ \hline & \text{blue diagonal} & & \end{array} \right] \end{array}$$

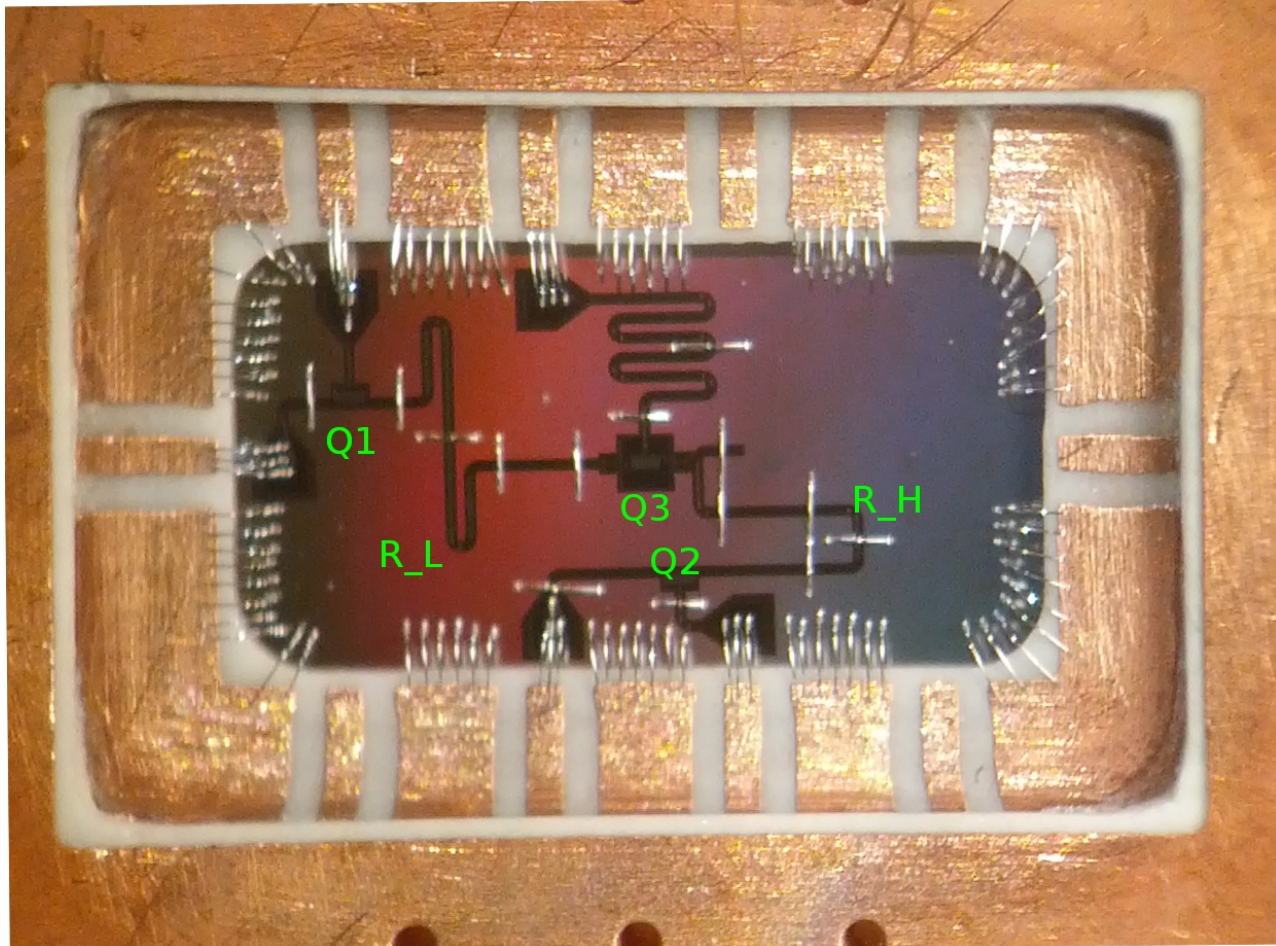
All without the systematic errors of previous procedures!

Experimental Procedure

- Repeat the following $\sim 16^n$ times
 - Pick a V
 - Run experiments of different lengths interleaving \mathcal{E} with random unitaries to compare $\mathcal{E} \circ \Lambda_U$ to V .
 - Run same experiments interleaving “null operation” with random unitaries to compare Λ_U to V .
- Use this data to reconstruct $\mathcal{E} \circ \Lambda_U$ and Λ_U
- Reconstruct \mathcal{E}

Process can take days...

Experimental Implementation



Questions?

BBN Quantum Information Group

Quantum Optics:

Journal of Modern Optics

Publication details, including instructions for authors and subscription information:

<http://www.tandfonline.com/loi/tmop20>

Approaching Helstrom limits to optical pulse-position demodulation using single photon detection and optical feedback

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^a Raytheon BBN Technologies, 10 Moulton Street, Cambridge, MA 02138, USA

^b National Institute of Information and Communications Technology, 4-2-1 Nukui-kitamachi,
.....

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Quantum Optics:

PHYSICAL REVIEW A **80**, 052310 (2009)

Gaussian-state quantum-illumination receivers for target detection

Saikat Guha¹ and Baris I. Erkmen²

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²*Jet Propulsion Laboratory, California Institute of Technology, Pasadena, California 91109, USA*

(Received 23 July 2009; published 10 November 2009)

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Super-Conducting Qubits:

Phys. Rev. Lett. 104, 163601 (2010) [4 pages]

Direct Observation of Coherent Population Trapping in a Superconducting Artificial Atom

Abstract

References

Citing Articles (20)

Download: PDF (680 kB) [Buy this article](#) Export: [BibTeX](#) or [EndNote \(RIS\)](#)

William R. Kelly^{*}, Zachary Dutton[†], John Schlafer, Bhaskar Mookerji, and Thomas A. Ohki[‡]
Raytheon BBN Technologies, Cambridge, Massachusetts 02138, USA

Jeffrey S. Kline and David P. Pappas
National Institute of Standards and Technology, Boulder, Colorado 80305, USA

Characterizing Quantum Operations

Shelby Kimmel

Joint work with Marcus Silva,
Raytheon BBN Technologies

Trading 1 Error for Another?

- Compare $\mathcal{E} \circ \Lambda_C$ to U for lots of U 's
 - Learn $(15^n + 1)$ of $(16^n - 4^n)$ parameters of $\mathcal{E} \circ \Lambda_C$
- Do same but remove \mathcal{E} : Compare Λ_C to U for lots of U 's
 - Learn $(15^n + 1)$ of $(16^n - 4^n)$ parameters of Λ_C

Clifford Twirl

- Cliffords are a set of unitaries
 - Apply a random Clifford \sim apply a random unitary
 - Can efficiently simulate classically

Clifford Twirl of \mathcal{E} :

Choose a random Clifford C

$$\llbracket \mathcal{E} \rrbracket = \text{---} \boxed{C} \text{---} \boxed{\mathcal{E}} \text{---} \boxed{C^\dagger} \text{---}$$

Clifford Twirl

Outcome is same for any initial state:

$$\rho \text{ --- } \boxed{[\mathcal{E}]} \text{ --- } p\rho + (1 - p)\rho_{\text{II}}$$

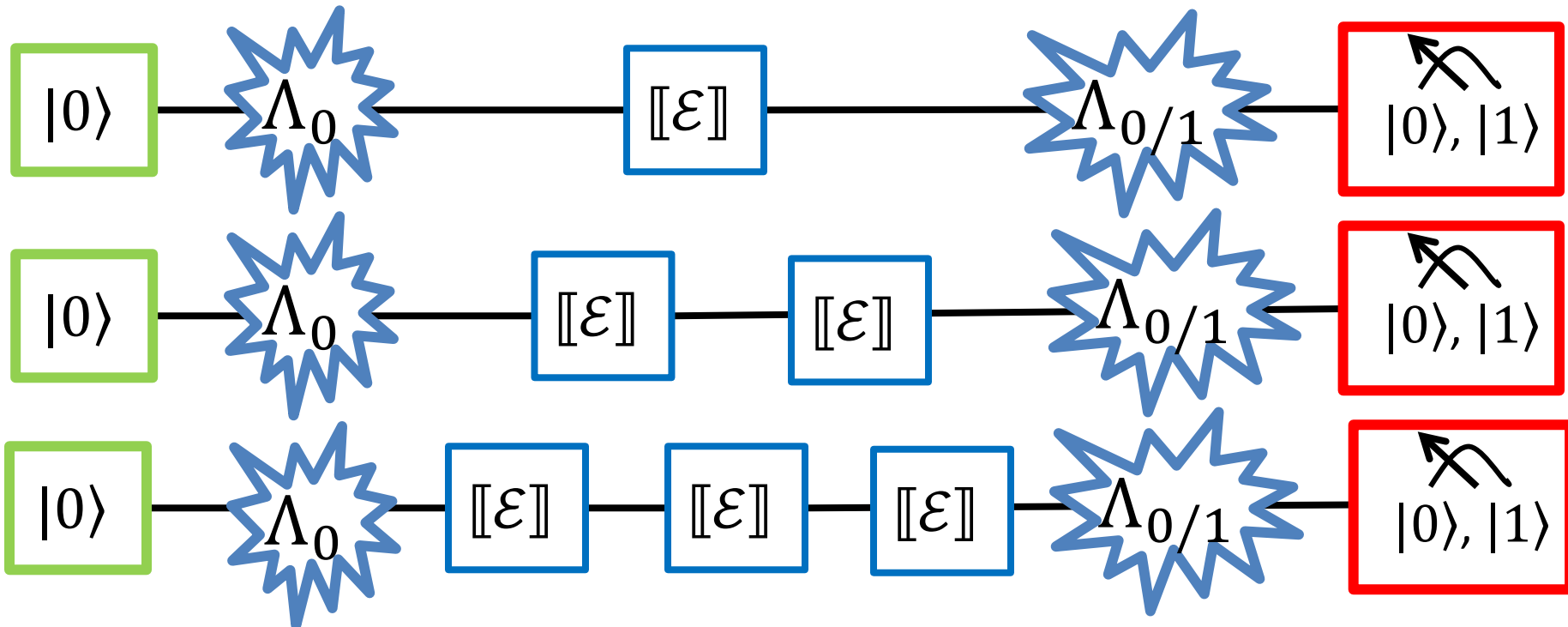
Composes Easily

$$\rho \text{ --- } \boxed{[\mathcal{E}]} \text{ --- } \boxed{[\mathcal{E}]} \text{ --- } \boxed{[\mathcal{E}]} \text{ --- } p^3\rho + (1 - p^3)\rho_{\text{II}}$$

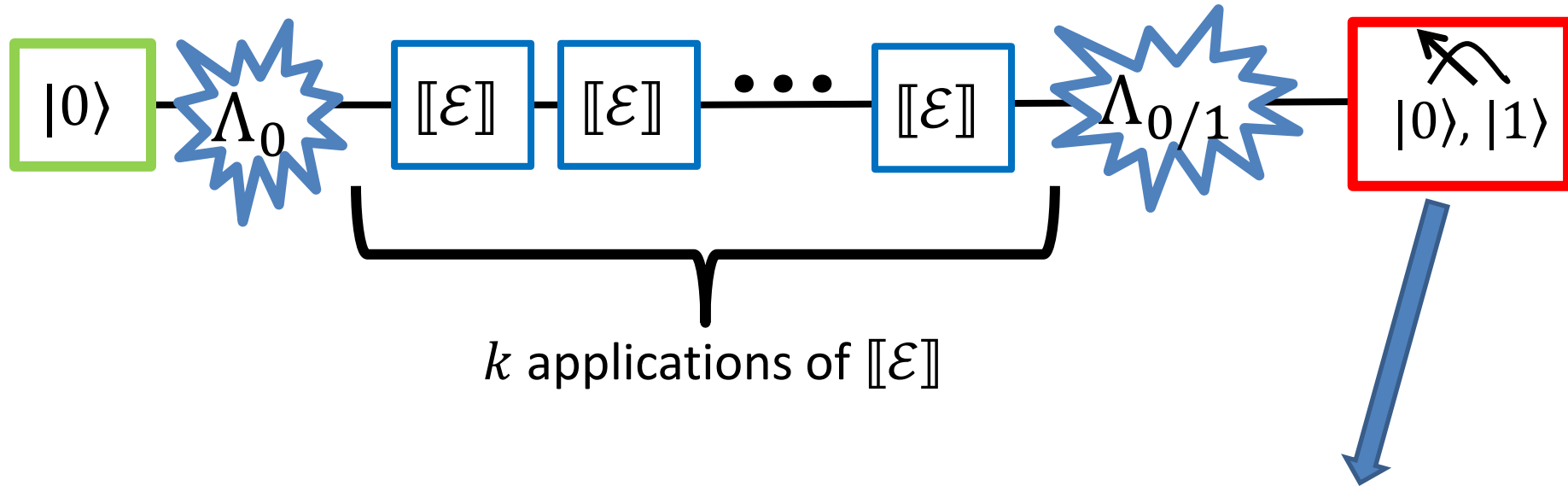
Repeated Applications

Choose a random Clifford C

$$[\mathcal{E}] = \text{---} \boxed{C} \text{---} \boxed{\varepsilon} \text{---} \boxed{C^\dagger} \text{---}$$



Repeated Applications



k applications of $[\mathcal{E}]$

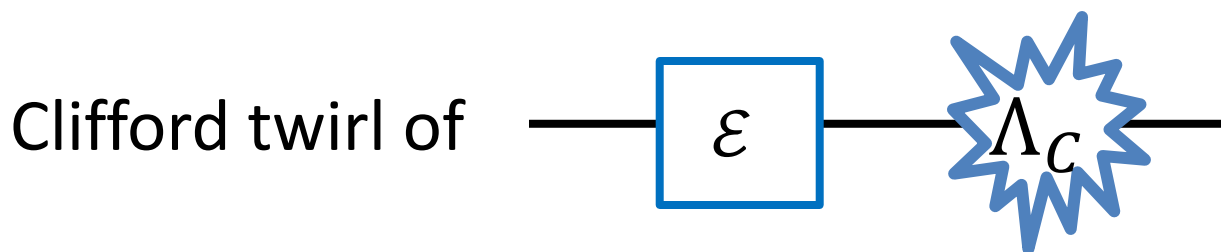
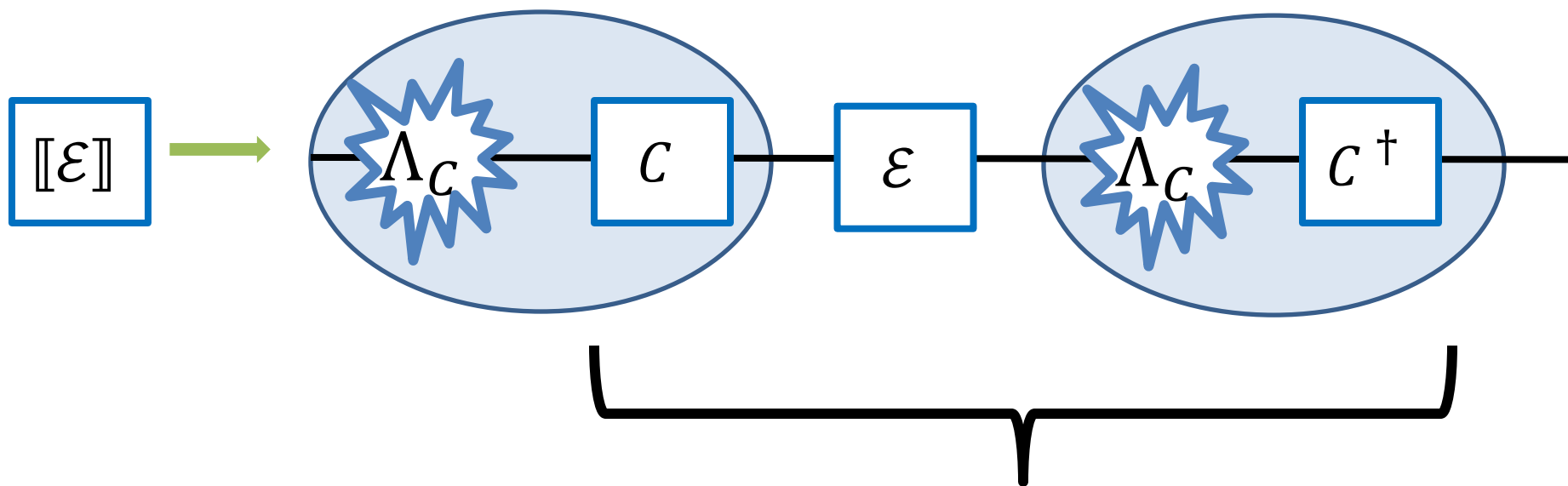
Measurement Outcomes:

$$Ap^k + B$$

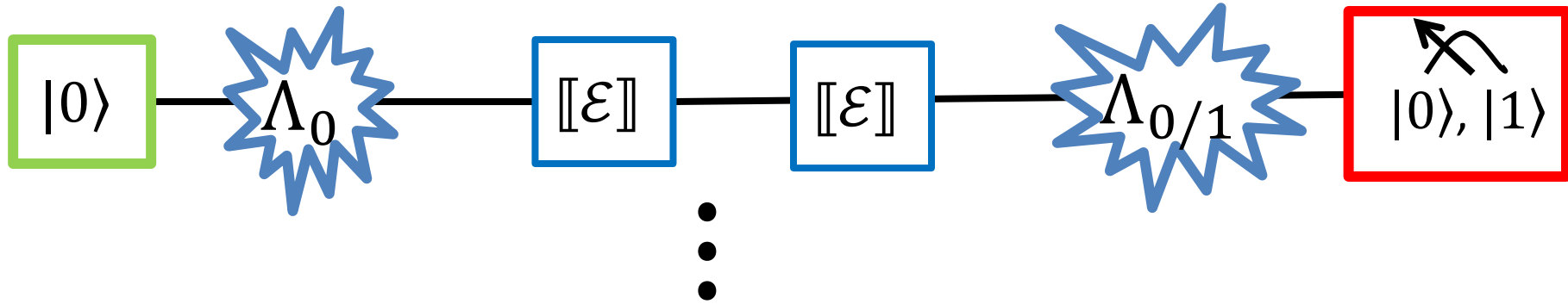
A, B depend on $\Lambda_0, \Lambda_{0/1}$

depends on closeness of \mathcal{E} to U

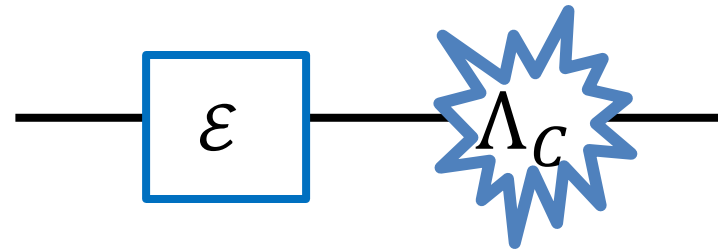
What about Errors!



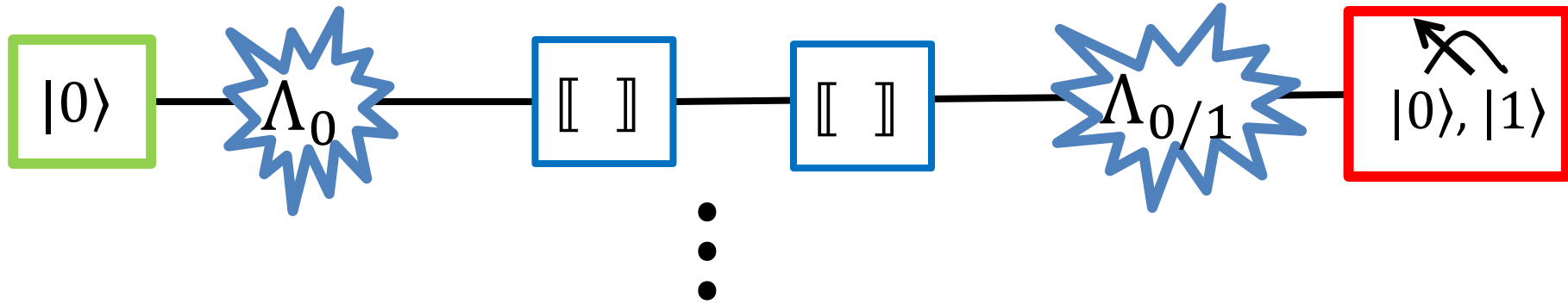
Procedure Step 1



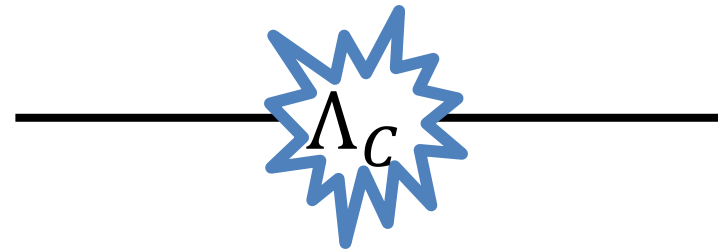
Learn p (probability of error) for



Procedure Step 2

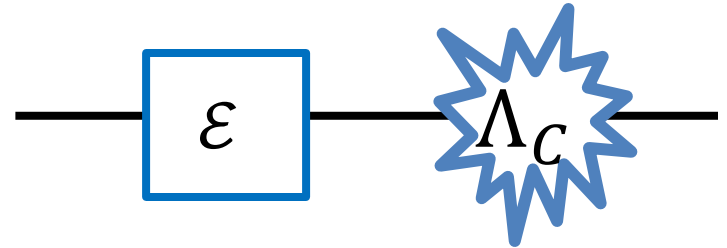


Learn p (probability of error) for



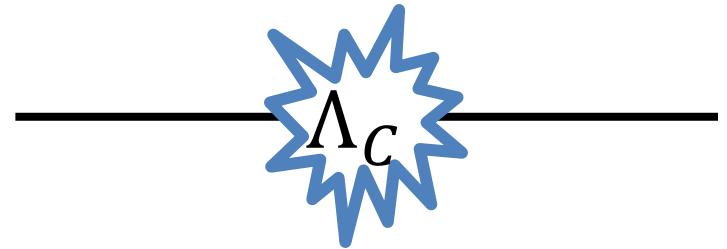
Procedure Step 3

Learn p (probability of error) for



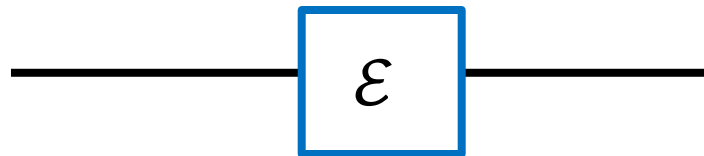
+

Learn p (probability of error) for



=

Bound on p for

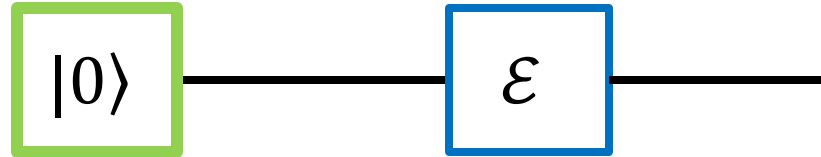


Need More Information!

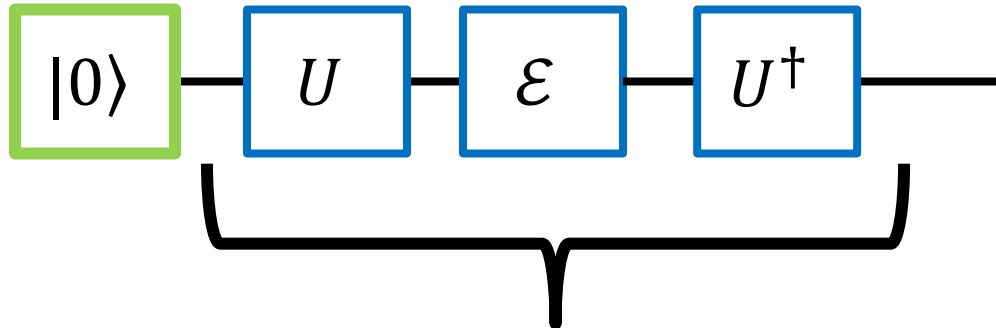
Only learn 1 parameter: p

Need $16^n - 4^n$ parameters to get full process

Twirling



What is action on an average state? Choose a random unitary U :



$$\boxed{[\varepsilon]} = \int dU \quad \boxed{U} \text{---} \boxed{\varepsilon} \text{---} \boxed{U^\dagger}$$

Twirling 2

$$\boxed{[\mathcal{E}]} = \int dU \quad \boxed{U} \text{---} \boxed{\varepsilon} \text{---} \boxed{U^\dagger}$$

$$\boxed{\rho} \text{---} \boxed{[\mathcal{E}]} \text{---} p\rho + (1-p)\frac{1}{d}\mathbb{I}$$

p depends on how similar is to identity

Randomized Benchmarking (imagine no Error)

Inserting Clifford

Unital Block

What about Error

Experimental Implementation

The Problem

