

A Multi-tool for your Quantum Algorithmic Toolbox

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Based on work with

- Stacey Jeffery: arXiv: 1704.00765 (Quantum vol 1 p 26)
- Michael Jarret, Stacey Jeffery, Alvaro Piedrafita, arXiv:1804.10591 (ESA 2018)
- Kai DeLorenzo, Teal Witter, arXiv:1904.05995 (TQC 2019)
- Bohua Zhan, Avinatan Hassidim, arXiv:1101.0796 (ITCS 2012)

Quantum Tools for Classical Problems

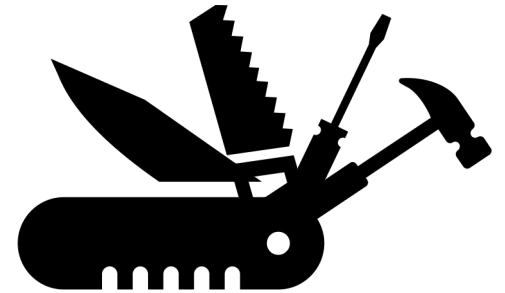
- Quantum Fourier Transform
- Grover Search
- Quantum Approximate Optimization Algorithm (QAOA)

Quantum Tools for Classical Problems

- Quantum Fourier Transform (good if highly structured)
- Grover Search (good if unstructured)
- Quantum Approximate Optimization Algorithm (QAOA)
(hard to analyze)

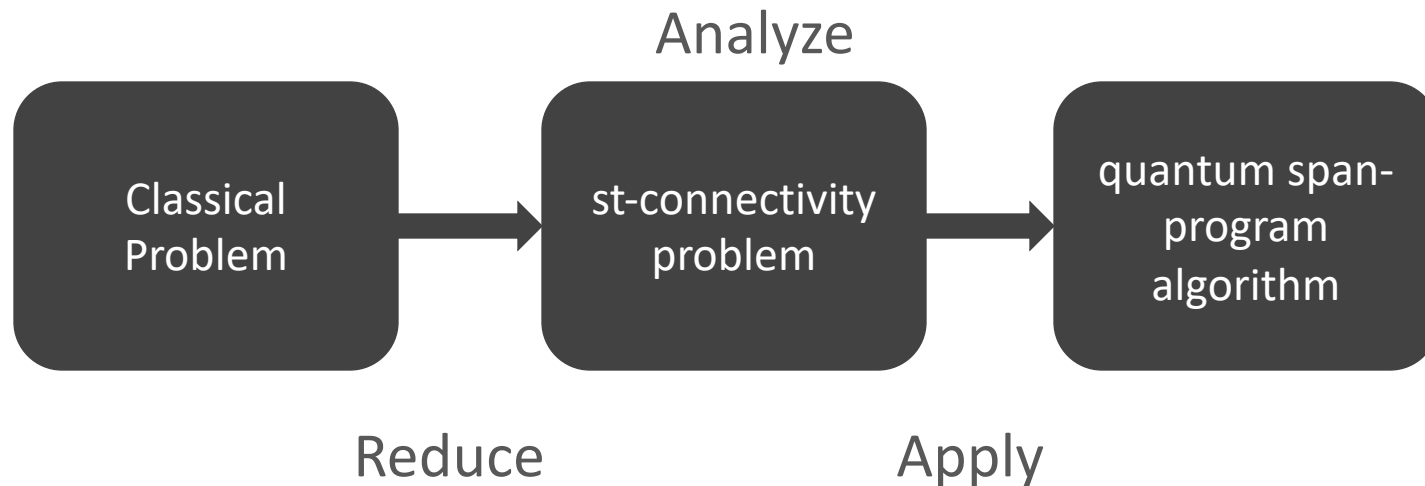
Multi-tool

- ✓ Structured
 - ✓ Unstructured
 - ✓ Easy creation
 - ✓ Easy, provable, non-quantum analysis
-
- ❖ Query model



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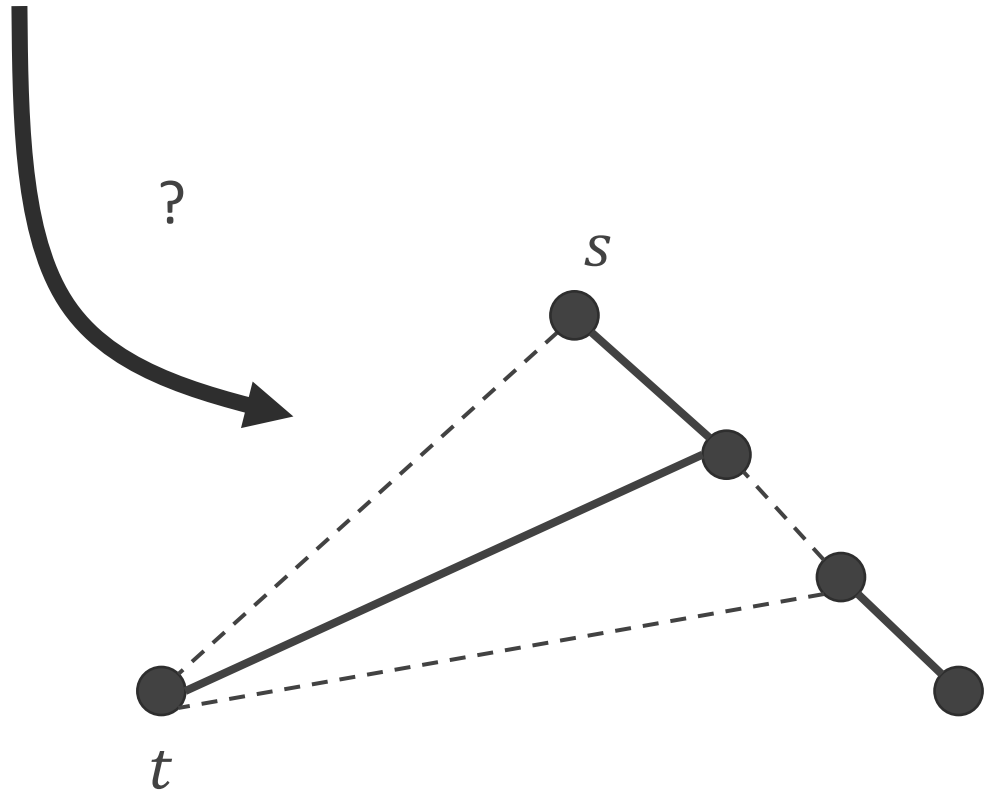


Reduction

Problem: Bit string $x = x_1x_2x_3$ contains a 1?

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Reduction (Decision Tree Approach)

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Idea: Turn classical algorithm into decision tree:

Algorithm:

- For each bit:
 - If 1: Output 1, end
- Output 0

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Start



●
Output 1

●
Output 0

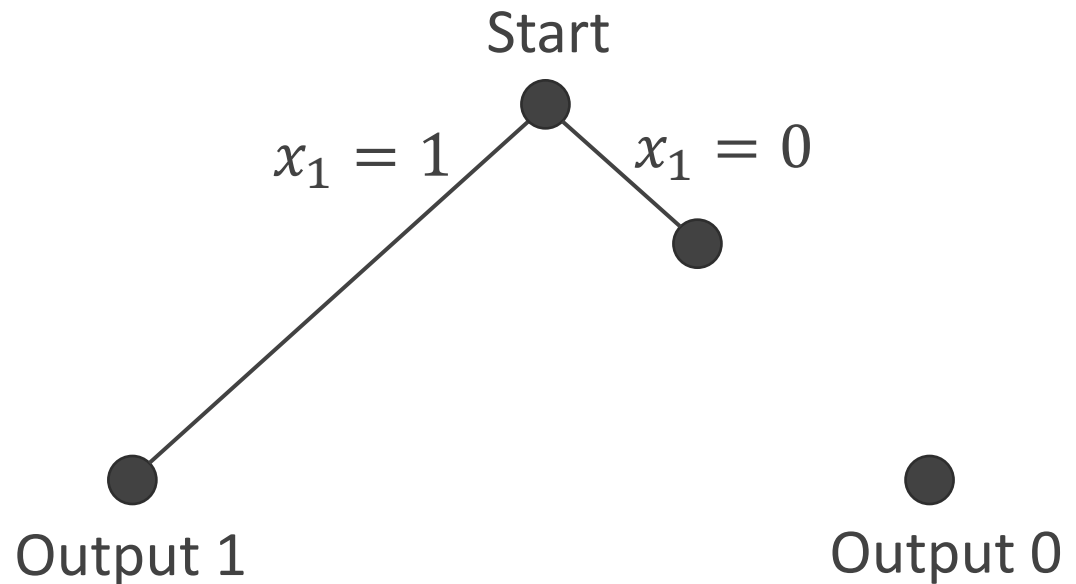
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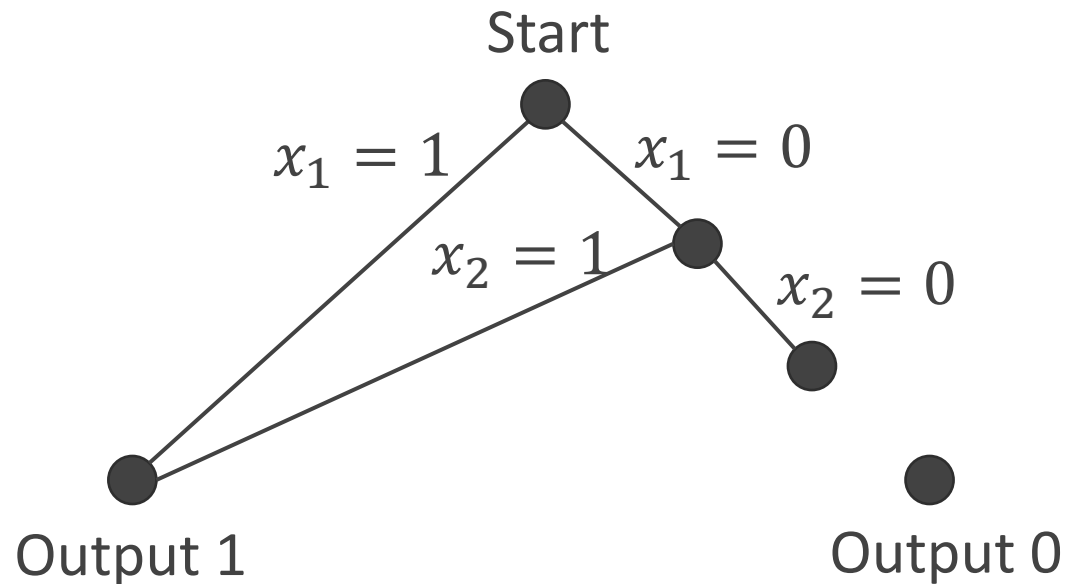
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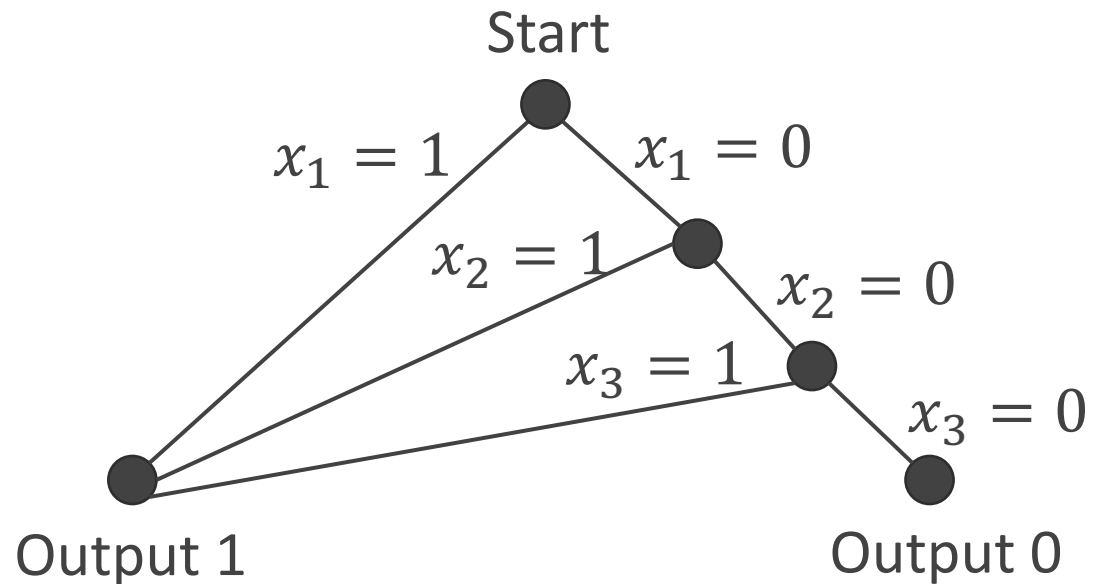
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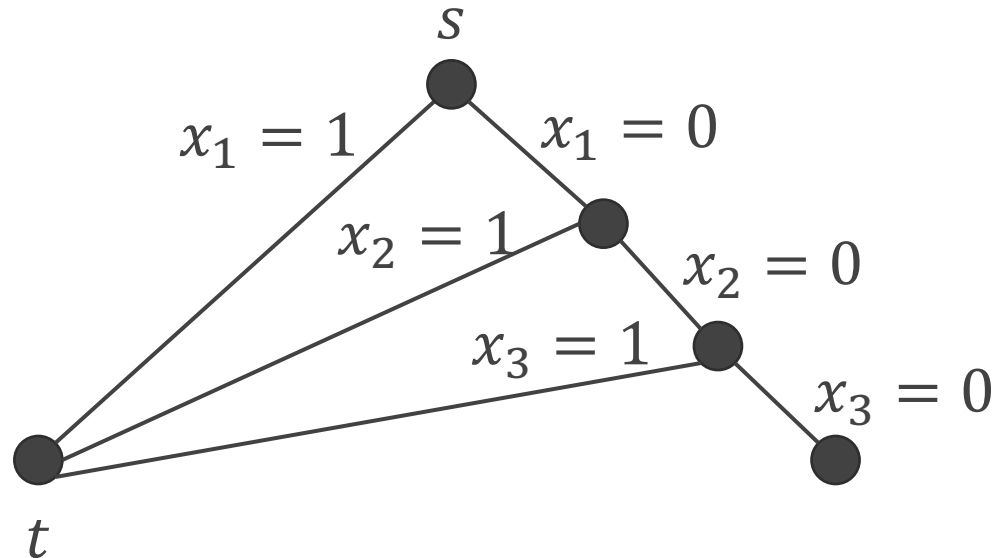
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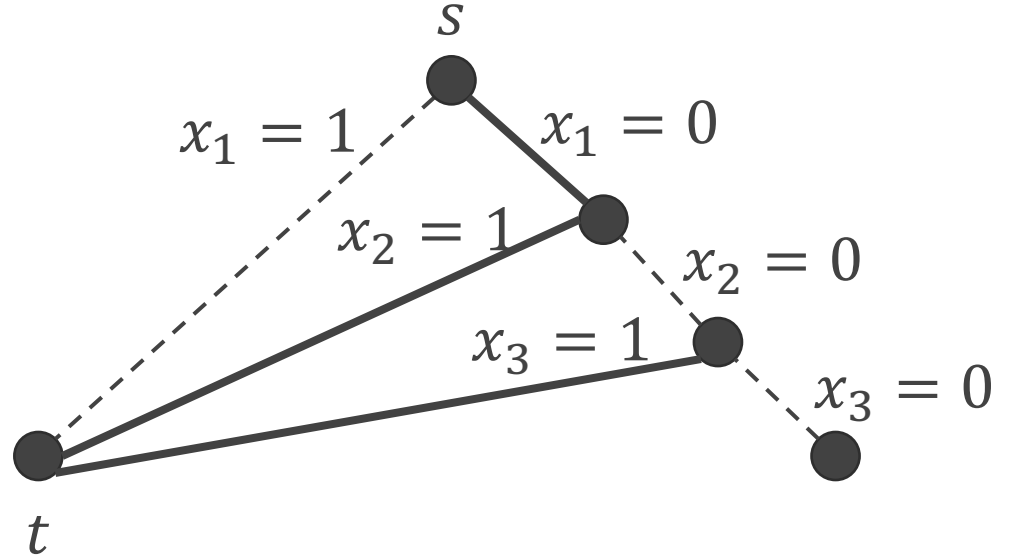


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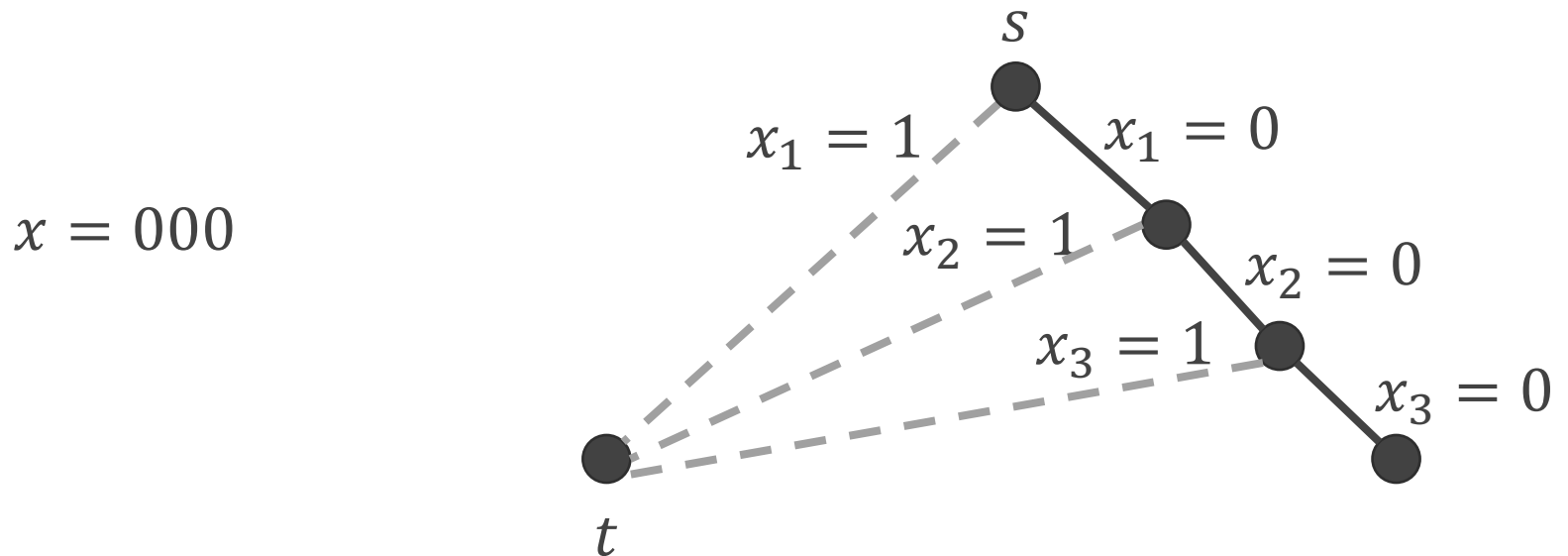
$x = 011$



Reduction (Decision Tree Approach)

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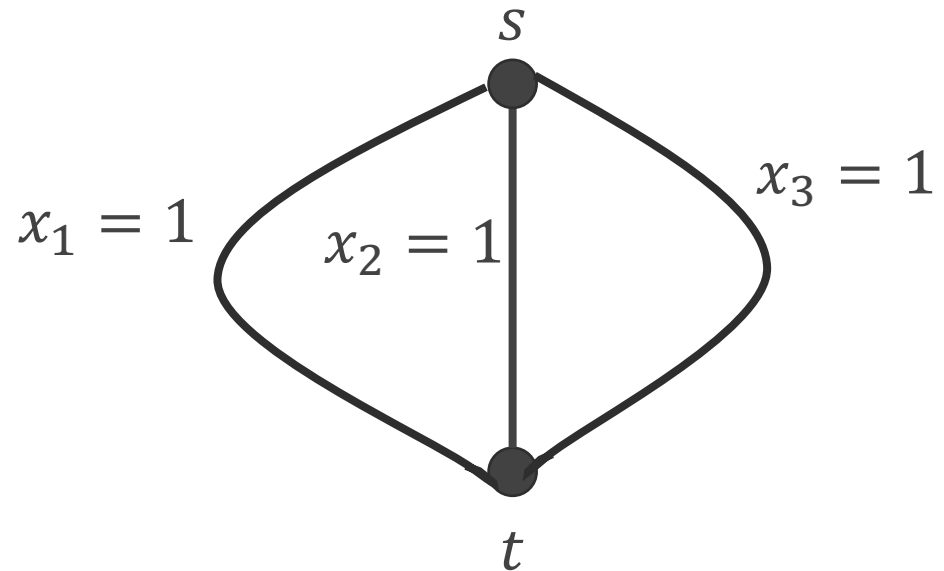
Idea: Turn classical algorithm into decision tree:



Reduction (Parallel Approach)

Problem: Bit string $x = x_1x_2x_3$ contains a 1?

Idea: Take advantage of parallel paths

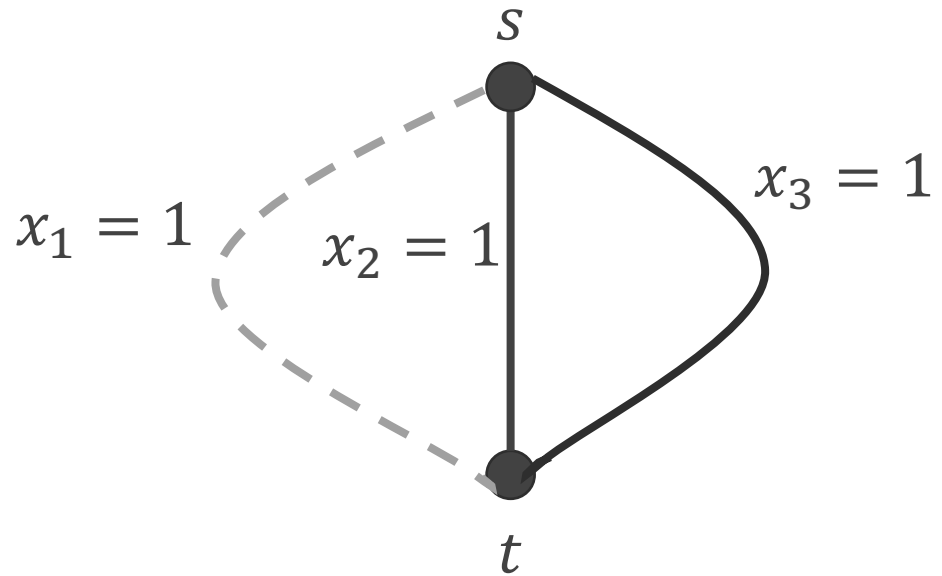


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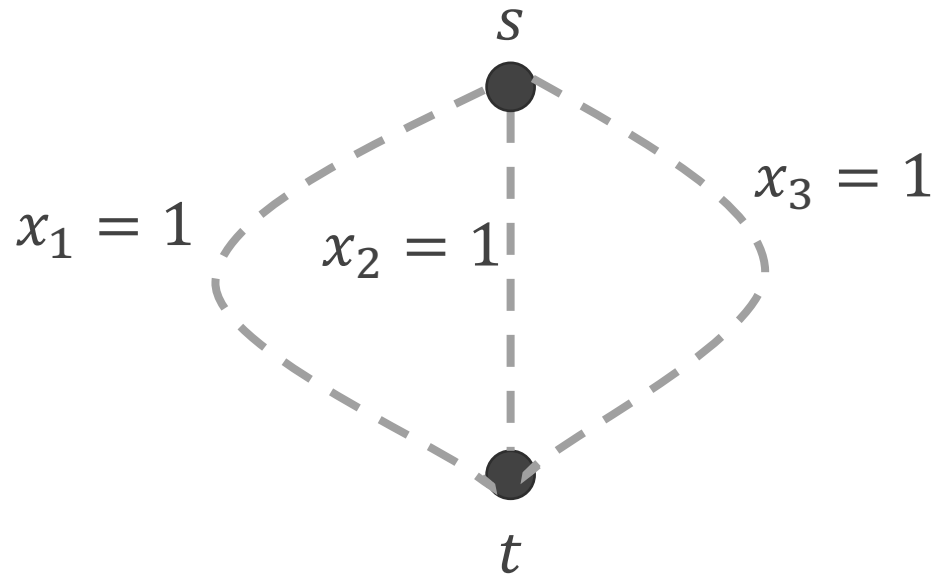


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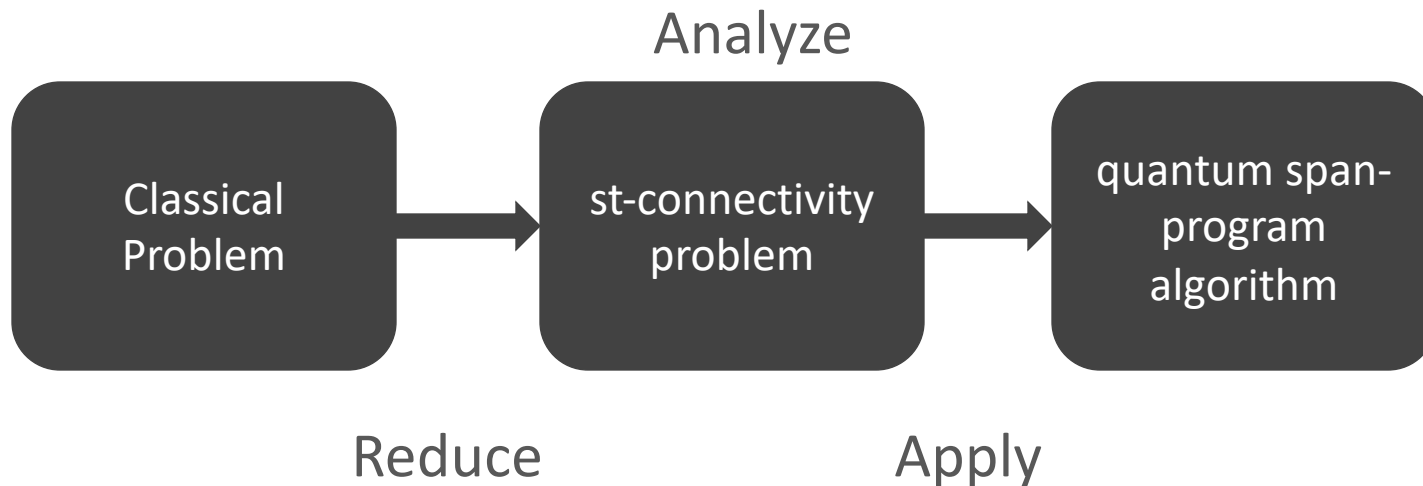
Idea: Take advantage of parallel paths

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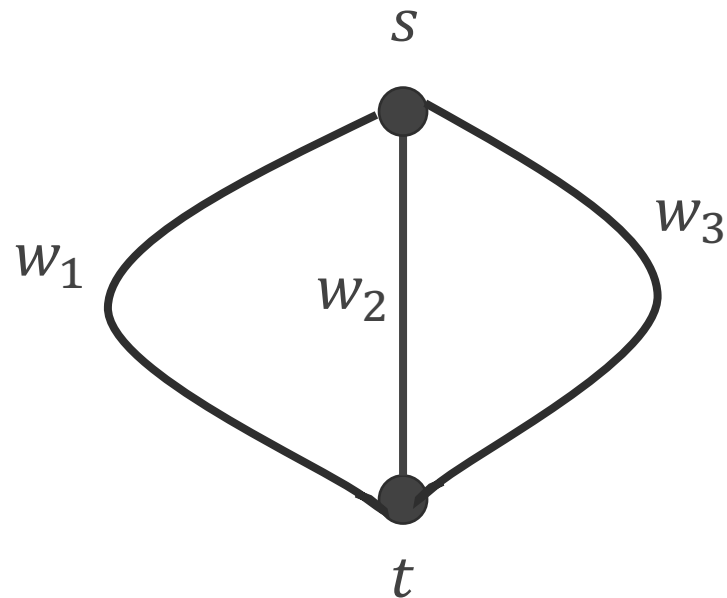
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Analysis

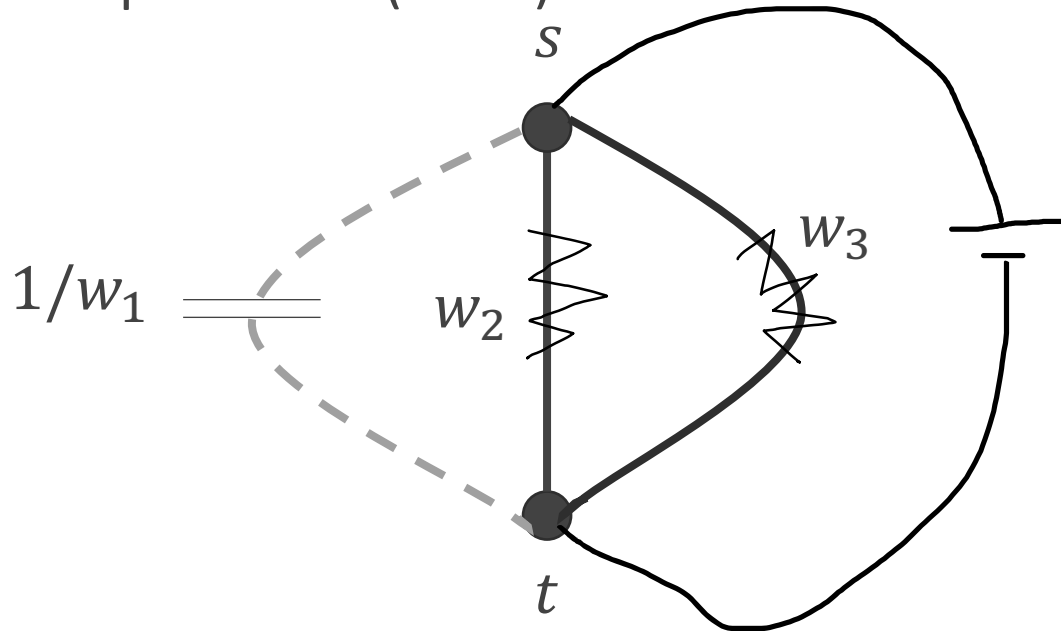
Assign a weight to each edge:



Analysis

For each input, can calculate

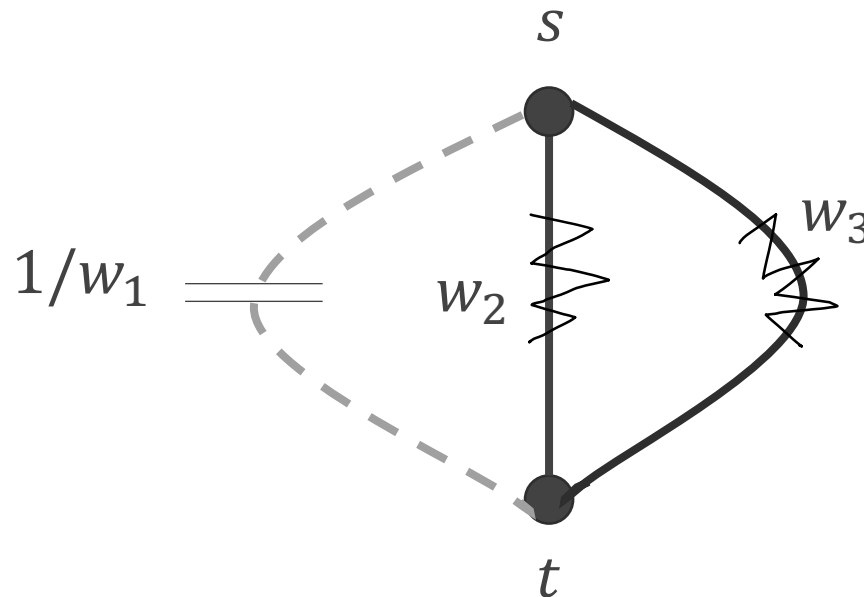
- effective resistance (if path)
- effective capacitance (if cut)



Analysis

Then Quantum Query Complexity is:

$$O(\sqrt{(\max C)(\max R)})$$

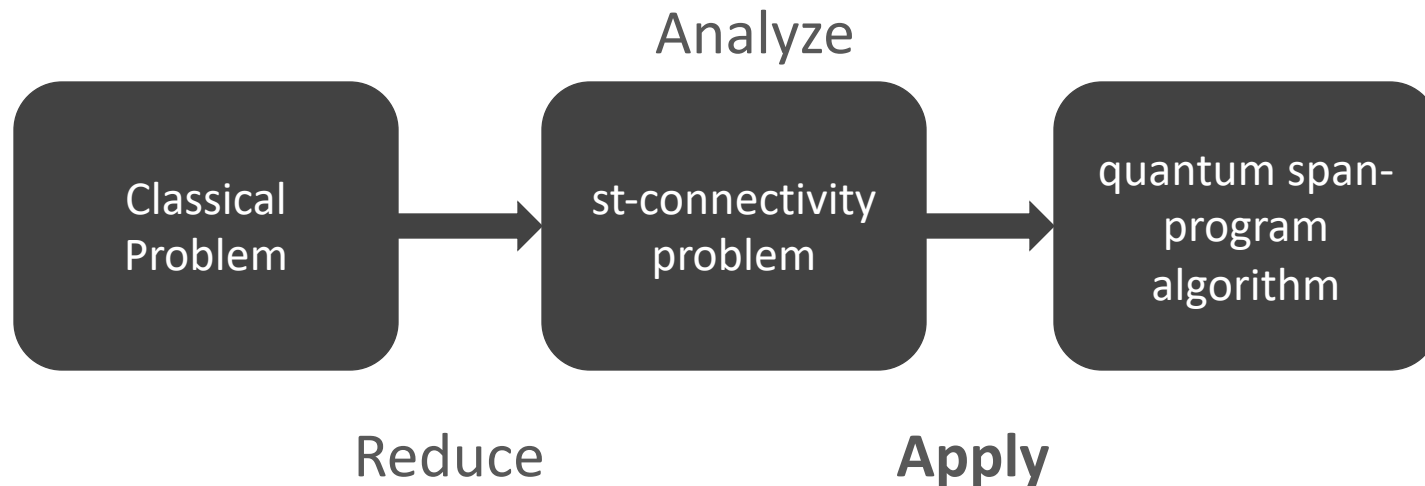


$\max R$: max effective resistance over connected instances

$\max C$: max effective capacitance over not connected instances

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Algorithm

Apply phase estimation to a unitary that is a product of two unitaries,

- One depends on input x
- One depends on underlying graph

Space complexity scales $\log(\text{no. edges, vertices})$ in graph.

[Belovs, Reichardt '12]

Performance

*There are alternative optimal algorithms for some problems

- Read-once Boolean formulas (query optimal) [JK]
- Total connectivity (query optimal) [JJKP]
- Cycle detection (query optimal) [DKW]
- Even length cycle detection [DKW]
- Bipartiteness (query optimal) [DKW]
- Directed st-connectivity (query optimal) (Beigi, Taghavi `19)
- Directed smallest cycle (query optimal) (Beigi, Taghavi `19)
- Topological sort (Beigi, Taghavi `19)
- Connected components (Beigi, Taghavi `19)
- Strongly connected components (Beigi, Taghavi `19)
- k-cycle at vertex v (Beigi, Taghavi `19)
- st-connectivity (query optimal) (Reichardt, Belovs `12)
- Maximum bipartite matching (Lin, Lin `16; Beigi and Taghavi `19)
- Maximum matching (**K**, Witter, in preparation)
- Super-polynomial speed-up for game evaluation [ZHK]

Summary

- A multi-tool for algorithm design that is accessible

Open Problems

- How to set weights?
- When is this approach optimal?

Funding:



Middlebury



ARO

People:



Stacey
Jeffery



Michael
Jarret



Alvaro
Piedrafita



Teal
Witter

Kai De
Lorenzo