Problems with Multiple Oracles

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What '



IT PROMISES TO SOLVE SOME OF HUMANITY'S MOST COMPLEX PROBLEMS. IT'S BACKED BY JEFF BEZOS, NASA AND THE CIA. EACH ONE COSTS \$10,000,000 AND OPERATES AT 459° BELOW ZERO. AND NOBODY KNOWS

HOW IT ACTUALLY WORKS





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What to do with a Quantum Computer?

• Let's try to solve something hard: 3-SAT

Traveling Salesman Problem [Karp '72] Tetris [Demaine et al '03]

Is a graph planar? [Grigoriev et al '07] Scheduling jobs [Ullman '75]

> Sudoku [Yato et al]

Graph Coloring [Karp '72]

3-SAT

• Goal: Want to satisfy a set of Boolean clauses, each with 3 variables.

$$(x_1 \vee \neg x_2 \vee x_3) \wedge (x_2 \vee x_3 \vee \neg x_4)$$

Each variable x_i can take value 0 or 1

V is logical OR:

$$0 \land 0 = 0$$

 $0 \land 1 = 1$
 $1 \land 0 = 1$
 $1 \land 1 = 1$
 \neg is logical NOT:
 $\neg 0 = 1$
 $\neg 0 = 1$
 $(\land 0 = 0)$
 $1 \land 0 = 0$
 $1 \land 0 = 0$
 $1 \land 0 = 0$
 $1 \land 1 = 1$

-1 = 0

3-SAT

Is there an assignment of the *n* variables $\{x_1, x_2, \dots, x_n\}$ such that $F(x_1, x_2, \dots, x_n) = 1$, where

$$F(x_1, x_2, \cdots, x_n) = (x_1 \lor \neg x_3 \lor x_7) \land (x_2 \lor x_3 \lor \neg x_6) \land (x_1 \lor x_7 \lor x_{10}) ..$$

~poly(n) clauses (e.g. Cn^2)

Algorithm for 3-SAT

$$F(x_1, x_2, \cdots, x_n) = (x_1 \vee \neg x_3 \vee x_7) \wedge (x_2 \vee x_3 \vee \neg x_6) \wedge (x_1 \vee x_7 \vee x_{10}) \dots$$

~poly(n) clauses (e.g. Cn^2)

Guess a satisfying assignment. Test if all clauses are satisfied
 Need to test ~2ⁿ possible inputs. With quantum computer can do in $\sqrt{2^n}$ steps

Outline

- Oracles and Oracle Models
- Related work
- Simple Example: Search with Multiple Oracles
- Open Problems and Directions for Future Work

Standard Oracle Model vs 3-SAT

3-SAT	Oracle Model
Given a function $F(x_1, x_2, \dots, x_n) = (x_1 \lor \neg x_3 \lor x_7)$	Initially function <i>F</i> is unknown
Determine if input x exists such that $F(x) = 1$	 Determine some property of <i>F</i> Does <i>x</i> exists such that <i>F</i>(<i>x</i>) = 1? Is <i>F</i> one-to-one?

Standard Oracle Model

Goal: Determine a property of a function F(x) for Boolean input $x = \{x_1, x_2, \dots, x_n\}$, given an oracle for F.



Only care about # of oracle calls (queries)

Multiple Oracles with Costs Model

Goal: Determine a property of a function F(x) for Boolean input $x = \{x_1, x_2, \dots, x_n\}$, given a set of oracles associated with functions $\{F_1, \dots, F_k\}$ which each have some information related to F

Care about total cost = $\sum_{i=1}^{k} q_i c_i$ where q_i is the # of times Oracle *i* is used

Related Work

- Ambainis '10: One oracle, but querying different *x* requires different amounts of time
 - E.g. To learn $f(00 \cdots 00)$ takes time 1, but to learn $f(11 \cdots 11)$ takes time 2
- Montanaro '09: Searching when given some additional information as to the location of the marked item.
 - E.g. Told that $f(00 \cdots 00) = 1$ is more likely than $f(11 \cdots 11)$
- Cerf et al. '00: Use multiple oracles to speed up evaluation of satisfiability problems.
 - No cost, No lower bounds, Need certain structure.

Searching with an Oracle

Goal: Determine x such that F(x) = 1. Can ask oracle, "Is the *i*th item the starred item?"

- Classically: $\sim N$ queries to oracle
- Quantumly: $\sim \sqrt{N}$ queries to oracle [Grover '97, Bennett et al. '97, Zalka '99]









Can ask \bigstar Oracle, "Is the *i*th item starred?" with cost c_{\star} Can ask Oracle, "Is the *i*th item striped?" with cost c_{\parallel}

C★>**C**∣

Promised: The starred item is also striped







Can ask \bigstar Oracle, "Is the *i*th item starred?" with cost c_{\star} Can ask () Oracle, "Is the *i*th item striped?" with cost c_{\parallel}

C_{*}>C|

Promised: The starred item is also striped



Sometimes best to check all N items using +Oracle

Otherwise can check all *N* items using cheaper oracle, but still need to use \bigstar Oracle for *M* items

Quantum Query Complexity Bounds



Size of Problem

• TOOL:

Need at least $\sim \sqrt{N}$ queries to quantum oracle to find one item out of N [Bennett et al '97]

* "at least $\sim \sqrt{N}$ " means at least $A\sqrt{N}$ for some constant A as N gets large. = $\Omega(\sqrt{N})$

• How much can **Oracle** help?



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- How much can **Oracle** help?
- At best, can narrow down to *M* items. But now searching for 1 among *M*, need at least $\sim \sqrt{M}$ queries to **Cracle TOOL**



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At best, can narrow down to *M* items. But now searching for 1 among *M*, need at least $\sim \sqrt{M}$ queries to **Cracle TOOL**



★ Oracle has cost c_{\star} , and need to use it $\sim \sqrt{M}$ times. Always will have a cost of $\sim c_{\star}\sqrt{M}$

- Need \bigstar Oracle $\sim \sqrt{M}$ times; How many times do we need \blacksquare Oracle?
- What if only needed to use () **Oracle** \sqrt{M} times? (For contradiction.)
 - This would be awesome!
 - Would always want to have an **()** Oracle, because it helps so much

Idea: Even if aren't given **()** Oracle, create it using + Oracle.











• What if only needed to use () Oracle \sqrt{M} times? (For contradiction.) Idea: Even if aren't given () Oracle, create it using \bigstar Oracle.

Can simulate **Oracle** using one (two) queries to **Cracle**!



Only need to use \bigstar Oracle $\sim \sqrt{M}$ times to simulate \blacksquare Oracle.

• What if only needed to use () Oracle \sqrt{M} times? (For contradiction.) Idea: Even if aren't given () Oracle, create it using \bigstar Oracle.

Can simulate \bigcirc **Oracle** using one (two) queries to \bigstar **Oracle**!





Previously showed need $\sim \sqrt{M}$ queries to \neq Oracle if have an \bigcirc Oracle.

• What if only needed to use () Oracle \sqrt{M} times? (For contradiction.) Idea: Even if aren't given () Oracle, create it using \bigstar Oracle.

Can simulate \bigcirc **Oracle** using one (two) queries to \bigstar **Oracle**!

- Only need to use \bigstar Oracle $\sim \sqrt{M}$ times to simulate \blacksquare Oracle.
- Previously showed need $\sim \sqrt{M}$ queries to \neq Oracle if have an \bigcirc Oracle.

Can find starred item using $\sim \sqrt{M}$ queries to \bigstar Oracle

• What if only needed to use () Oracle \sqrt{M} times? (For contradiction.) Idea: Even if aren't given () Oracle, create it using \bigstar Oracle.

Can simulate \bigcirc **Oracle** using one (two) queries to \bigstar **Oracle**!

Only need to use \bigstar Oracle $\sim \sqrt{M}$ times to simulate \blacksquare Oracle.

M queries to 🗡 🛛

Previously showed need $\sim \sqrt{M}$ queries to \Rightarrow Oracle if have an \bigcirc Oracle.

• Using this argument:



Always need to use either \bigstar Oracle or () Oracle at least $\sim \sqrt{N}$ times.



Algorithm for Searching with Multiple Oracles

Amplitude Amplification

Quantum Query Complexity Bounds



Size of Problem

Algorithm for 3-SAT

$$F(x_1, x_2, \cdots, x_n) = (x_1 \vee \neg x_3 \vee x_7) \wedge (x_2 \vee x_3 \vee \neg x_6) \wedge (x_1 \vee x_7 \vee x_{10}) \dots$$

~poly(n) clauses (e.g. Cn^2)

- Guess a satisfying assignment. Test if all clauses are satisfied EXPENSIVE
 Need to test ~2ⁿ possible inputs. With quantum computer can do in √2ⁿ steps
- Guess a satisfying assignment. Test if ~log(n) clauses are satisfied CHEAP
 Defines a subset of possible solutions, including the true satisfying assignment, if it exists
 What is M?

Directions for Future Work

Create tight bounds for searching with multiple oracles

 Adversary Bound/Span programs
 Geometric picture

- Can we create a general framework for understanding oracles with costs, in the way that the adversary bound is a framework for understanding standard oracle problems
- Many quantum oracle problems does it make sense to add multiple oracles to these problems?

Classical Algorithm

- 1. Choose item at random and test if striped using Oracle 2
- 2. If it is striped, test if starred using Oracle 1

Worst case cost:

 $c_1(M-1) + c_2N$